High-Speed Digital and Mixed-Signal Components for X– and KU–Band Direct Digital Synthesizers in Indium Phosphide DHBT Technology

Steven Eugene Turner

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HIGH–SPEED DIGITAL AND MIXED–SIGNAL COMPONENTS
FOR X– AND K_U–BAND DIRECT DIGITAL SYNTHESIZERS IN
INDIUM PHOSPHIDE DHBT TECHNOLOGY

By

Steven Eugene Turner
B.S. University of Maine, 2001
M.S. University of Maine, 2003

A THESIS
Submitted in Partial Fulfillment of the
Requirements for the Degree of
Doctor of Philosophy
(in Electrical Engineering)

The Graduate School
The University of Maine
May, 2006

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FOR X– AND KU–BAND DIRECT DIGITAL SYNTHESIZERS IN
INDIUM PHOSPHIDE DHBT TECHNOLOGY

By Steven Eugene Turner

Thesis Advisor: Dr. David E. Kotecki

An Abstract of the Thesis Presented
in Partial Fulfillment of the Requirements for the
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May, 2006

Recently reported double heterojunction bipolar transistor (DHBT) devices manufactured in Indium Phosphide (InP) technology with $f_t$ and $f_{max}$ both over 300 GHz enable advanced high-speed digital and mixed-signal circuits. In this thesis, the use of InP DHBT devices for high-speed accumulator circuits and X– and Ku–band direct digital synthesizer (DDS) circuits are investigated. At these frequencies, new technological challenges in the design of digital and mixed-signal circuits arise in areas including power consumption and clock distribution. This thesis addresses the speed/power trade-offs in high-speed accumulator designs, the design of DDS circuits, and clock distribution simulation. The results of six accumulator circuits and two DDS circuits are reported as part of this thesis. The fastest 4-bit accumulator at a 41 GHz clock rate is reported, as well as the fastest DDS circuits operating at 13 GHz and 32 GHz clock rates. The 13 GHz DDS has a worst case spurious-free dynamic range (SFDR) of 26.67 dBc and consumes 5.42 W of power, while the 32 GHz DDS has a worst case SFDR of 21.56 dBc and consumes 9.45 W of power. In addition to the circuit designs, a methodology for simulating electrically long clock interconnects and a new figure of merit for comparing DDS designs are developed.
ACKNOWLEDGMENTS

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<th>Abbreviation</th>
<th>Description</th>
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<tbody>
<tr>
<td>CML</td>
<td>current mode logic</td>
</tr>
<tr>
<td>CMOS</td>
<td>complementary metal-oxide-semiconductor</td>
</tr>
<tr>
<td>DAC</td>
<td>digital to analog converter</td>
</tr>
<tr>
<td>DDS</td>
<td>direct digital synthesizer</td>
</tr>
<tr>
<td>DHBT</td>
<td>double heterojunction bipolar transistor</td>
</tr>
<tr>
<td>DTL</td>
<td>diode-transistor logic</td>
</tr>
<tr>
<td>DUT</td>
<td>device under test</td>
</tr>
<tr>
<td>ECL</td>
<td>emitter coupled logic</td>
</tr>
<tr>
<td>( f_0 )</td>
<td>fundamental output frequency</td>
</tr>
<tr>
<td>( f_{\text{clk}} )</td>
<td>clock frequency</td>
</tr>
<tr>
<td>FCW</td>
<td>frequency control word</td>
</tr>
<tr>
<td>FPGA</td>
<td>field programmable gate array</td>
</tr>
<tr>
<td>GSGSG</td>
<td>ground-signal-ground-signal-ground</td>
</tr>
<tr>
<td>HBT</td>
<td>heterojunction bipolar transistor</td>
</tr>
<tr>
<td>InP</td>
<td>Indium Phosphide</td>
</tr>
<tr>
<td>LPF</td>
<td>low pass filter</td>
</tr>
<tr>
<td>LSB</td>
<td>least significant bit</td>
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<tr>
<td>LUT</td>
<td>look up table</td>
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<tr>
<td>MIM</td>
<td>metal-insulator-metal</td>
</tr>
<tr>
<td>MSB</td>
<td>most significant bit</td>
</tr>
<tr>
<td>MSFOM</td>
<td>mixed-signal figure of merit</td>
</tr>
<tr>
<td>RAM</td>
<td>random access memory</td>
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<tr>
<td>RTL</td>
<td>resistor-transistor logic</td>
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<tr>
<td>SFDR</td>
<td>spurious-free dynamic range</td>
</tr>
<tr>
<td>SHBT</td>
<td>single heterojunction bipolar transistor</td>
</tr>
<tr>
<td>TFAST</td>
<td>Technology for Agile Digitally Synthesized Transmitters</td>
</tr>
<tr>
<td>TTL</td>
<td>transistor-transistor logic</td>
</tr>
<tr>
<td>VBIC</td>
<td>Vertical Bipolar Inter-Company</td>
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</table>
CHAPTER 1

Introduction

Continuing advances in heterojunction bipolar transistor (HBT) technology are expanding the range of applications for HBT devices. Recently, HBTs have been reported in InP with $f_t$ and $f_{max}$ both over 300 GHz [1]. As the switching time of HBT devices decreases, they become useful in high-speed digital and mixed-signal circuits. The best method for taking advantage of the inherent speed of HBTs for digital and mixed-signal circuits is by using emitter coupled logic (ECL). Using the standard emitter coupled pair shown in Figure 1.1 as a current steering switch building block, complex logic functions can be constructed. ECL is fast because currents are switched over small signal swings, typically 200 mV to 300 mV, and node capacitances are charged and discharged over small ranges [2]. Beyond high-speed operation, ECL is advantageous because it incorporates differential signals, making it less susceptible to noise [3].

ECL has been around for many years, and it is used in circuits where high-speed operation is the most critical factor. Examples of applications include microprocessors [4], random access memories (RAMs) [5], radars [6], field programmable gate arrays (FPGAs) [7], and frequency synthesizers [8]. For the past couple of decades, most of the work done in high-speed microprocessors and RAMs has moved away from ECL and into complementary metal-oxide-semiconductor (CMOS) technologies. In these applications, the lower power consumption of CMOS technologies outweighs the extra speed provided by ECL and HBTs. However, ECL is still used for niche applications, such as state of the art advances in radars, FPGAs, and frequency synthesizers. In these applications, higher frequencies are usually desired at the expense of higher power consumption. Despite the longevity of ECL, the ever increasing clock frequencies in high-speed digital and mixed-signal applications bring about many challenges in areas such as power consumption and clock distribution, that must be dealt with.
1.1 Major Design Challenges for High-Speed Digital and Mixed-Signal HBT Circuitry

Two of the most difficult challenges that must be dealt with in the design of high-speed digital and mixed-signal HBT circuitry are power consumption and clock distribution. The two are related, and both become increasingly difficult to address as circuit complexity and clock speeds increase. For logic gates operating at the maximum switching speed in the Vitesse VIP-2 process [1], ECL circuits must have current densities around 5 mA/µm². This leads to basic gates (inverters, NANDs, etc.) requiring around 15 mA to 20 mA, or 54 mW to 72 mW of power at a 3.6 V supply voltage. Even at the lower figure, a circuit with 20 basic gates consumes more than 1 W of power. Non-trivial systems for applications such as radars, frequency synthesizers, and FPGAs require gate counts in the hundreds or thousands. For these high power gates, that translates into tens to hundreds of Watts of power.
1.1.1 Power Consumption Issues

In circuits and systems, high power consumption is problematic for a number of reasons. First, high power consumption leads to increased die and device temperatures. Direct measurements of the temperature at the die and wafer level is often difficult to do without expensive equipment. An indirect measure of temperature for wafers and dies is power density, which is defined as the amount of power consumption per unit area. Typically, a conservative maximum power density limit of 100 W/cm$^2$ is cited [9] for the ability to remove heat from a die in either HBT or CMOS technologies. As the power density and temperature increase, the performance of the HBT devices and circuits decrease. As a result, steps must be taken to reduce the power density or to remove heat from the die. While gates can be spaced out on the die to decrease junction temperatures, large dies increase manufacturing costs, since more wafers must be processed to produce the same number of chips. When there are large spaces between gates, physically long lines must be used for gate connections. These long lines are prone to the effects of increased interconnect capacitance. Additionally, if the lines are electrically long (see the Section 1.1.2), microwave design techniques must be employed, further complicating the design process. While heat sinks, liquid cooling, and fans can be employed to compensate for elevated die temperatures, they can add extra power consumption and physical weight at the system level, decreasing portability.

Even if the heat related problems are ignored, high power consumption limits the potential applications for high-speed circuits and systems. While it is possible to supply power on the order of tens or hundreds of Watts in a vehicle or a fixed location, massive amounts of power are not feasible for portable or hand-held systems. Since many of the potential applications for high-speed digital and mixed-signal circuits lie within the field of communications, it is advantageous to keep power consumption low so that hand-held communication devices are viable.
1.1.2 Clock Distribution Issues

Systems with any sort of complexity require clock distribution to many registers and latches on locations throughout the chip. On highly integrated circuits with high-speed clock signals, the distribution lengths of the clock interconnects can become electrically long, so that they can no longer be modelled as lumped element interconnects. One estimate for when interconnects begin to be electrically long and the lumped element model begins to be insufficient is when their length is greater than \( \frac{1}{6} \) of the effective length of a rising edge of the signal \([10]\). This can be expressed as

\[
    l = \frac{T_r}{6D}. \tag{1.1}
\]

With an 11 ps rise time \((T_r)\) for a 30 GHz clock signal and a delay \((D)\) of 0.0067 ps/\(\mu\)m in InP, interconnects longer than 300 \(\mu\)m can be considered to be electrically long. Note that the delay figure is an estimate for the top metal layer of the four metal layer stack up in the Vitesse VIP-2 technology \([1]\). This layer is typically used for clock distribution in the circuits in this thesis. Signal interconnects other than clock lines can be electrically long.

The clock signal must also arrive at all of the distribution points with reasonably matched phase, or low clock skew, in order to reduce setup and hold violations in the latches \([11]\). While all of the clock interconnects can be designed to have the approximately same length to minimize clock skew, the Spectre simulation toolkit does not directly deal with the effects of electrically long interconnects correctly. Instead, the parasitic extraction of the Spectre toolkit incorrectly models the electrically long lines with lumped element models that have large lumped parasitic capacitors connected. This sort of model ignores the distributed resistance, capacitance, and inductance that would be present in a transmission line model.
In order to work around this issue, separate microwave models of the clock tree are constructed and simulated with the Spectre toolkit to determine proper line matching and to investigate the behavior of long line effects. Problems such as excessive gain or attenuation at certain frequencies can be addressed by adding terminations or source resistors to the long lines. Once it is determined that the interconnects are well matched and well behaved in a microwave sense, they are assumed to be safe in the circuit simulations, even though no direct models can be simulated for the long interconnects in the Spectre simulator. While this method gives reasonable results, there are many potential pitfalls in using two separate simulation environments.

Clock distribution problems tend to require solutions that increase power consumption. This includes the addition of buffers in the distribution path or high-power emitter followers for better drive capability. The buffers can be used for driving a large number of clock branches in a distribution tree, or for recovering weak signals. High-power emitter followers are used for driving the clock to compensate for parasitics and transmission line effects. As system complexity increases, the clock distribution can be a large proportion (40% or more [12]) of the total power for the circuit, and all of the problems related to high power consumption are further aggravated.

### 1.2 Power/Speed Trade-Off

High power consumption is generally necessary to achieve the high switching speeds necessary for high-speed circuits. Consequently, lower power designs typically sacrifice speed for power. Figure 1.2 shows a typical $f_t$ curve, which is a general figure of merit in HBT designs. The $f_t$ curve shows that the switching speed is reduced for current densities below the peak $f_t$. Since the current is proportional to the power, lower power leads to lower switching speed. However, it is still possible to maintain high clock rates while lowering power by reducing supply voltages, using alternative architectures, or optimizing timing paths.
1.2.1 Timing Path Optimization

One method for decreasing the power consumption of a circuit is to take advantage of the excess timing margin in gates that are part of non-critical paths. The excess timing margin can exist when some logic gates are faster than needed to complete an operation within a clock cycle. Since the speed of the gates can be reduced by reducing the current density and the corresponding power, the gates with excess timing margin are effectively wasting power. For gates in non-critical paths, the power and resulting speed can be reduced to some extent before the overall circuit performance is impacted. Thus, it is possible to reduce the power with no impact on the overall circuit performance. In some cases, gates can be operated at less than half power without a negative impact on the overall circuit performance.

Power can be optimized by analyzing the timing characteristics of all of the data paths and reducing the power in gates which have excess timing margin. Careful analysis and power optimization of the system timing paths can lead to significant power
savings. It is not necessary to specifically set the power for each instance of a gate. Instead, a few versions of each type of gate with different power and speed characteristics can be created and used for optimization purposes. This scaled back method of optimization can allow for reduction in power that can be implemented relatively rapidly in terms of engineering time.

1.2.2 Architecture Modifications

Additional reduction in power can be achieved through modifications to the circuit architecture or topology. Typically, high-speed digital systems require many stages of pipelining in order to achieve high clock rates. Pipeline stages increase power consumption by adding registers and extra clock circuitry that is needed to drive the registers. In systems with the highest clock rates, the combination of registers and clock circuitry easily accounts for a majority of the total system power. Any architecture modifications that can reduce the number of registers will have a large impact on reducing the power consumption.

1.2.3 Supply Voltage Reduction

Power consumption can also be reduced by decreasing the supply voltage. However, this involves the elimination of logic levels in the gates. For example, in order to implement a 3-input ECL AND gate, at least 3 NPNs must be stacked as shown in Figure 1.3. In the AND gate in Figure 1.3, the inputs and outputs are differential. The output differential pair is Zp/Zn, and the input differential pairs are Ap/An, Bp/Bn, and Cp/Cn. This requires a voltage supply of ~4.5 V in order to accommodate the minimum voltage swing (200 mV to 300 mV), 4 diode drops (0.8 V/diode), and the voltage required for the current source (1.0 V). By removing a level from the logic and using a 2-input AND gate, a diode drop can be removed from the supply voltage, reducing it to about 3.7 V. This helps to decrease the power consumption, since many of the gates,
such as registers, are 2-input gates. It is important to note that while using a reduced supply voltage decreases the power consumed by individual gates, it may potentially increase the overall power if more gates are needed to implement a logic function. For example, a 3-input gate must be implemented as two 2-input gates. However, this is compensated for in systems with many pipeline stages, where registers comprise a large portion of the total number of gates, and the registers do not need any additional current to operate with a lower supply voltage. In these systems, the reduction in the power consumption of the registers is much greater than the increase in the power consumption in the logic, so the overall power is reduced.

1.3 Motivation for Project

This thesis is supported in part by Defense Advanced Research Projects Agency (DARPA), which is pursuing high performance mixed-signal circuits through the Technology for Frequency Agile Digitally Synthesized Transmitters (TFAST) program [13].
Table 1.1: Goals of the TFAST program.

<table>
<thead>
<tr>
<th></th>
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<tbody>
<tr>
<td>Emitter width</td>
<td>0.25 μm</td>
<td>0.15 μm</td>
</tr>
<tr>
<td>Current Density</td>
<td>500 kA/cm²</td>
<td>1000 kA/cm²</td>
</tr>
<tr>
<td>$f_t/f_{max}$ (with $BV_{CEO} = 4.5$ V)</td>
<td>350 GHz / 400 GHz</td>
<td>500 GHz / 500 GHz</td>
</tr>
<tr>
<td>Static Divider Clock Speed</td>
<td>150 GHz</td>
<td>200 GHz</td>
</tr>
<tr>
<td>Transistor Count</td>
<td>1 k</td>
<td>20 k</td>
</tr>
</tbody>
</table>

Direct digital synthesizer (DDS) circuits are of particular interest for applications including radars, communications, and electronic warfare [9]. SiGe has some advantages over InP, since it is generally more manufacturable, has better lithography, has smaller devices, and it can realize high levels of circuit complexity, than InP devices. The TFAST program, however, is using InP HBTs, because they have much higher breakdown voltages and slightly better $f_t$ and $f_{max}$, resulting in a better mixed-signal performance than InP. The TFAST program is also promoting the development of InP HBT technologies to continue to increase their speed advantages while making them in highly manufacturable processes that are commercially viable for large transistor count circuits. The goals of the TFAST program [14] are shown in Table 1.1. The Phase I goals were reached in mid-2004.

This thesis focuses on the high-speed digital and mixed-signal circuit design for direct digital synthesizers implemented in InP HBT technology. Of particular interest are DDS circuits that can synthesize frequencies over X-band (8 GHz to 12 GHz) and K_u-band (12 GHz to 18 GHz) frequencies. In order to maximize the potential applications of the high-speed DDS circuitry, it is also desirable to achieve the lowest power consumption possible. Since power and speed are two conflicting metrics, the trade-offs between them will be explored. Generally, increasing the power of HBT circuits increases the speed, but the increased power reaches a point of diminishing returns where large increases in power yield small benefits for speed performance. Eventually, further power increases actually have a negative impact on the speed after the peak point of
the $f_t$ curve is reached, as shown in Figure 1.2. This thesis investigates the trade-offs involved in maximizing clock speed while minimizing power through the use of various circuit architectures and design techniques. While there are numerous speed/power design points depending on the specific application, this thesis will provide some guidance for choosing an architecture and using design techniques to yield the lowest possible power at a particular clock frequency in InP HBT technology. This thesis will also develop a figure of merit for comparing DDS designs.

1.4 Review of Previously Reported Work on Digital and Mixed-Signal HBT Circuits

The DDS is a mixed-signal system, since it contains both digital and analog components. For InP and other III-V HBTs, one way to measure the relative performance of mixed-signal circuits is through the mixed-signal figure of merit (MSFOM), defined by [14]:

$$MSFOM = \frac{J_C \cdot BV_{CEO}}{C_{CB} \cdot \Delta V_{LOGIC}}, \quad (1.2)$$

where $J_C$ is the current density, $BV_{CEO}$ is the breakdown voltage, $C_{CB}$ is the extrinsic capacitance between the base and collector, and $\Delta V_{LOGIC}$ is the voltage swing. In Equation 1.2, the bottom term ($C_{CB} \cdot \Delta V_{LOGIC}$) represents a time constant from parasitics and voltage swing. As this time constant increases, the MSFOM is reduced. Likewise, increasing $J_C$ tends to overcome this time constant. Higher breakdown voltages improve the MSFOM, since a higher $BV_{CEO}$ contributes to higher linearity [15] in digital to analog converters (DACs), which are critical components of mixed-signal circuits. In a purely digital InP HBT system, the major delay term [16] is analogous to the inverse of the MSFOM without $BV_{CEO}$.

From Equation 1.2, it is apparent that factors other than the $f_t$ and $f_{max}$ of a technology play a role in determining its relative usefulness for mixed-signal circuits.
While the MSFOM provides a more complete metric, not all of the components necessary for computing the MSFOM are always available in the literature. For digital systems, one figure of commonly reported merit is the maximum clock frequency for a static frequency divider. While this metric neglects the $BV_{CEO}$ term needed for the MSFOM, it gives a relative estimate for the $J_C/(C_{CB} \cdot \Delta V_{LOGIC})$ term of the equation. When compared to the $f_t$ and $f_{max}$, the maximum frequency of the static frequency divider provides an estimate of what level of performance degradation arises from layout parasitics in simple digital circuits.

While useful as an estimate, the maximum static frequency divider is not always a realistic metric. Reported static frequency dividers are usually designed to reach a maximum frequency and are not necessarily practical in more complex circuits. These frequency dividers are often operated at very high current densities which are only beneficial to circuit performance when the circuit is continuously oscillating. The dividers sometimes use peaking inductors to improve high frequency response \[17]-[19] but in more complex circuits the peaking inductors are often impractical due their layout size and the interconnect routing limitations they introduce.

Reported static dividers implemented in III-V technologies (including InP) are summarized in Table \[1.2\] and static dividers implemented Si-bipolar and SiGe technologies are summarized in Table \[1.3\]. The most recently reported dividers in both tables operate at approximately $f_t/2$. Variations in earlier reported static dividers tended to be worse than $f_t/2$ for III-V technologies and better than $f_t/2$ for Si-bipolar/SiGe technologies. It has been noted that divider performance is limited by factors other than $f_t$, such as base resistance ($R_b$) \[20\], collector capacitance ($C_{JC}$) \[21\], and interconnect parasitics \[22\]. A graph of the maximum frequency of Si-bipolar/SiGe and III-V static frequency dividers versus the year reported is shown in Figure \[1.4\].

For a complete comparison of MSFOM and high-speed mixed-signal circuits, it is more appropriate to study reported DDS circuits. However, many of the reported
<table>
<thead>
<tr>
<th>Year Reported</th>
<th>Frequency (GHz)</th>
<th>$f_t$ (GHz)</th>
<th>Company or Group</th>
<th>Reference</th>
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<tbody>
<tr>
<td>1989</td>
<td>34.8</td>
<td>68</td>
<td>NTT</td>
<td>[23]</td>
</tr>
<tr>
<td>1991</td>
<td>36</td>
<td>110</td>
<td>HRL</td>
<td>[24]</td>
</tr>
<tr>
<td>1992</td>
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<td>HRL</td>
<td>[25]</td>
</tr>
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<td>1998</td>
<td>52.9</td>
<td>134</td>
<td>HRL</td>
<td>[26]</td>
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<td>1999</td>
<td>66</td>
<td>164</td>
<td>UCSB</td>
<td>[17]</td>
</tr>
<tr>
<td>1999</td>
<td>69</td>
<td>195</td>
<td>TRW</td>
<td>[18]</td>
</tr>
<tr>
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<td>HRL</td>
<td>[22]</td>
</tr>
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<td>165</td>
<td>UCSB</td>
<td>[27]</td>
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<td>205</td>
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<td>[28]</td>
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<tr>
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<td>135$^b$</td>
<td>HRL</td>
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<td>405$^c$</td>
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<td>2004</td>
<td>152</td>
<td>301</td>
<td>UCSB/Rockwell</td>
<td>[16]</td>
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<tr>
<td>2004</td>
<td>152</td>
<td>300$^d$</td>
<td>BAE/Vitesse</td>
<td>[1]</td>
</tr>
</tbody>
</table>

$^a$ Measurements were limited by available test equipment.

$^b$ The $f_{\text{max}}$ of this process is over 300 GHz.

$^c$ The $f_{\text{max}}$ of this process is 370 GHz.

$^d$ $f_t$ is reported in the range of 300 GHz to 350 GHz.

Table 1.2: Reported static frequency divider circuits implemented in III-V processes.

<table>
<thead>
<tr>
<th>Year Reported</th>
<th>Frequency (GHz)</th>
<th>$f_t$ (GHz)</th>
<th>Company or Group</th>
<th>Reference</th>
</tr>
</thead>
<tbody>
<tr>
<td>1989</td>
<td>12.5</td>
<td>30</td>
<td>Fujitsu</td>
<td>[31]</td>
</tr>
<tr>
<td>1991</td>
<td>21</td>
<td>40</td>
<td>NEC</td>
<td>[32]</td>
</tr>
<tr>
<td>1993</td>
<td>25</td>
<td>45</td>
<td>Siemens</td>
<td>[33]</td>
</tr>
<tr>
<td>1995</td>
<td>30</td>
<td>NR$^a$</td>
<td>Siemens</td>
<td>[34]</td>
</tr>
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<td>1996</td>
<td>35</td>
<td>50</td>
<td>Siemens</td>
<td>[35]</td>
</tr>
<tr>
<td>1997</td>
<td>42</td>
<td>68</td>
<td>Siemens</td>
<td>[36]</td>
</tr>
<tr>
<td>2000</td>
<td>67</td>
<td>122</td>
<td>Hitachi</td>
<td>[37]</td>
</tr>
<tr>
<td>2001</td>
<td>71</td>
<td>123</td>
<td>Hitachi</td>
<td>[38]</td>
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<td>2003</td>
<td>86</td>
<td>200</td>
<td>Infineon</td>
<td>[39]</td>
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<tr>
<td>2003</td>
<td>96</td>
<td>210</td>
<td>IBM</td>
<td>[40]</td>
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</table>

$^a$ $f_t$ not reported.

Table 1.3: Reported static frequency divider circuits implemented in Si-bipolar and SiGe processes.
DDS circuits are implemented in CMOS technologies, and they operate at low frequencies (less than 200 MHz) to either illustrate new design techniques or to achieve high spurious-free dynamic range (SFDR). Also, technologies can not be compared since many of the high-speed DDS circuits are reported in data sheets or manuals that do not specify the technology used. An idea of the current state of the art for high-speed DDS systems can be derived from Table 1.4. Both commercial and research DDS circuits are included in Table 1.4. For commercial DDS circuits, AD9858 [41] provides an $f_{out\text{max}}$ of 400 MHz and an SFDR of 50 dBc while consuming 2 W of power. The research examples from A. Gutierrez-Aitken et al. at TRW [8] and K. Elliott at HRL [42] are closest to the type of work presented in this thesis. The TRW DDS was reported in an InP HBT technology in 2001 with an $f_{out\text{max}}$ of 4.56 GHz, an SFDR of 30 dBc, and a power consumption of 15 W. The HRL DDS was reported in an InP HBT technology in
2005. It had an $f_{out\text{max}}$ of 5.5 GHz, an SFDR of 30 dBc, and a power consumption of 8 W.

### 1.5 Thesis Organization

Chapter 2 discusses a general overview of the Vitesse InP technology, design tools, and ECL logic design techniques used in the thesis. Circuit design elements that are common to many of the designs in the thesis, such as current sources, voltage swings, and voltage levels are described in detail. An overview of the methods used for simulation of transmission lines with the simulation toolkit is also given.

High-speed accumulators are discussed in Chapter 3. The chapter is divided into four major sections. In Section 3.1, the basics of adder and accumulator design are discussed. Section 3.2 discusses four high-speed 4-bit accumulator designs. These accumulators operate up to a 41 GHz clock frequency. In Section 3.3, two 8-bit accumulator designs are discussed. These designs sacrifice some high-speed performance to achieve lower power consumption. The designs are also integrated into DDS circuits.
Finally, Section 3.4.1 discusses other techniques for power consumption reduction, namely triple-tail circuits and resistor-only current sources.

Two full DDS circuits that use the accumulators from Section 3.3 are described in Chapter 4. The DDSV1 design described in Section 4.4 operates up to a maximum clock frequency of 32 GHz and consumes 9.45 W of power, while the DDSV2 design described in Section 4.5 operates up to a maximum clock frequency of 13 GHz and consumes 5.42 W of power. A new figure of merit for comparing the performance of DDS designs is also introduced that takes the SFDR and frequency resolution of a design into account.

The thesis is concluded by a summary of accomplishments and suggestions for future work that are presented in Chapter 5.
CHAPTER 2
HBT Circuit Design

This chapter provides background information that is relevant to the circuits designed in this thesis. An overview of the Vitesse VIP-2 InP process is given in Section 2.1. This technology is used for manufacturing the circuits designed in this thesis. An overview of the simulation tools and methodology is presented in Section 2.2. The final sections of this chapter cover emitter coupled logic, voltage swing, and current sources.

2.1 Vitesse VIP-2 Process

Details of the Vitesse VIP-2 process used in this thesis were first reported in [1] and described in further detail in [48]. The VIP-2 process uses double heterojunction bipolar transistors (DHBTs). The DHBT devices have an n-InP emitter, a p-InGaAs base, and an n-InP collector. A diagram of the DHBT device is shown in Figure 2.1. The process has 4 levels of aluminum metal interconnect, thin-film resistors, and metal-insulator-metal (MIM) capacitors. There are no PNP devices in the process, only NPN transistors. A stack-up of the process is shown in Figure 2.2. The process is self-aligned, using dielectric and metal spacers. The use of self-alignment and the elimination of the traditional lift-off process results in a highly manufacturable process. In the VIP-2 process, \( f_t \) and \( f_{\text{max}} \) are both over 300 GHz, while \( BV_{CEO} \) is over 4 V.

Although single heterojunction bipolar transistors (SHBTs) are generally faster than DHBTs implemented in the same materials, the DHBTs have some advantages over SHBTs [49], such as higher breakdown voltages and better thermal characteristics. These advantages make the DHBT devices better suited for highly integrated digital and mixed-signals circuits.
Figure 2.1: Diagram of the self-aligned DHBT device [1] from the Vitesse VIP-2 InP process.

Figure 2.2: Stack-up of the Vitesse VIP-2 InP DHBT process [1] with 4 aluminum metal interconnect levels, thin-film resistors, and MIM capacitors.
The higher breakdown voltage in the Vitesse VIP-2 DHBT devices is due to the fact that the p-InGaAs to n-InP base to collector junction has a larger energy gap than the p-InGaAs to n-InGaAs base to collector junction in an equivalent SHBT device. A higher breakdown voltage is advantageous for use in mixed-signal circuitry, since a higher $BV_{CEO}$ allows for a larger DAC output range that leads to improved DAC linearity [15]. That is one of the reasons that $BV_{CEO}$ is included in the MSFOM outlined in Equation [1.2] in Chapter [1] because DAC circuits are essential to mixed-signal circuitry and a higher breakdown voltage will improve the MSFOM if other factors are constant. The higher breakdown voltage is also one of the main reasons that InP DHBTs are used instead of SiGe devices, since the VIP-2 process $BV_{CEO}$ is over 4 V, compared to below 2 V for SiGe [9].

The DHBT devices are also advantageous compared to similar SHBT devices, because the InP collector in a DHBT device is more thermally conductive than the InGaAs collector in a SHBT device [49]. High thermal conductivity is essential, particularly in densely packed mixed-signal circuits with thousands of transistors. The heat dissipated from the high performance devices needs to be moved away from the devices efficiently in order to minimize self-heating and reduce the the operating temperature. This ensures that the DHBT devices perform optimally.

The Vitesse VIP-2 process facilitates high-speed digital and mixed-signal circuits. Reported digital circuits include a 152 GHz static frequency divider [1] and a 41 GHz 4-bit accumulator [50]. For mixed-signal circuits, a 50 GHz variable gain amplifier [48] and DDS circuits operating up to a 32 GHz clock frequency [51] have been reported. The 41 GHz accumulator and 32 GHz DDS will be discussed in further detail in later sections of this thesis.
2.2 Circuit Simulation

The Cadence design environment is used for schematic capture and layout. Unless otherwise noted, Cadence Spectre is used for circuit simulations. The Vitesse VIP-2 transistors are modelled by the Vertical Bipolar Inter-Company (VBIC) model [52]. In most cases, transient time-domain simulations are used for simulation of the digital and mixed-signal circuits. When other types of simulations are used, it is noted. When simulations are run to determine the maximum clock frequency of a circuit, the input clock frequency is usually stepped in 1 GHz increments until the circuit fails. Failures in the circuits generally are caused by timing margin failures in the digital portions of the circuitry. Thus, the failures aren’t evidenced by a reduced differential, but instead by an incorrect output caused by an incorrect internal digital sequence. For example, in divide by two test circuits, a failed output will not be a true divide by two signal. Instead, it will “slip,” missing cycles when driven above its the maximum operating frequency.

For increased correlation between simulation and measured results, the simulation results almost always include parasitic extracted capacitances from layout. Unless otherwise noted, the simulation results presented in this thesis will include extracted parasitic capacitance. Schematic-only simulations without parasitics are also run for all of the circuits, but this intermediate step is usually excluded from this thesis, because the extracted parasitic capacitors have a significant impact on the designs. The schematic-only simulation results give guidance on design parameters, but the parasitic capacitors that are introduced from the layout greatly impact on the performance of the circuit and can not be ignored. In many of the designs, the layout is iterated and adjusted to reduce the parasitic capacitance if possible. Another option is to increase the drive of the emitter followers that are loaded by the large parasitics. Compared to experimental results, the simulations with parasitic capacitances extracted from layout generally provide a reasonable match.
2.3 Clock Distribution Simulations

The parasitic extraction of this technology is limited since it only includes capacitance. It does not include the parasitic resistance or inductance that is also inherent in fabricated circuits. Most interconnects are electrically short, so the parasitic capacitances are the dominant factor affecting circuit performance and they can be treated as a lumped parameter. As mentioned in Section 1.1.2 the clock interconnections can be electrically long and will not be modelled correctly by the transient time-domain simulations in Spectre using just lumped parameters.

For these situations, the clock tree is modelled by microstrip transmission lines that include the inherent distributed R, L, and C, along with the circuits that load the clock distribution interconnects. Models for transmission lines are included in Spectre, but they do not have corresponding models in layout. As a result, a clock tree schematic of the clock distribution is constructed and simulated with a frequency-domain AC simulation. In the clock tree schematics, the transmission line parameters, such as the characteristic impedance and line length, are extracted from the layout by hand, because the design kit does not include an automated method for extraction of these parameters. The clock tree schematic is simplified from the overall circuit schematic and includes the circuits that drive the clock and the circuits that are loads on the clock. All other circuitry is omitted from the simplified schematic. An example of a clock tree schematic for a 4-bit accumulator is shown in Figure 2.3.

Unlike most of the other simulations in this thesis that use time-domain analysis, Spectre frequency-domain AC simulations are used for the clock tree simulations. These simulations aid in pinpointing frequencies where the clock signals feeding the registers are potentially over-damped or under-damped. Based on the initial simulations, adjustments are made to the clock tree, such as adding buffers and source resistors, to compensate and adjust the gain to appropriate levels over the targeted clock frequency range. Once the clock gain is within appropriate levels, it is assumed that the clock
Figure 2.3: Example clock tree schematic for a 4-bit accumulator. The transmission line parameters are determined from the layout. The clock drivers and clock loads are included for the simulations.

circuitry will perform correctly, and the full circuit simulations are then run using time-domain transient analysis.

An example of this type of simulation output is shown in Figure 2.4. In these simulations, the clock signal magnitude is normalized so that clock levels above 0 dB are strong enough to drive registers. Based on the initial simulations that did not include series resistors in the clock distribution path, the worst case register inputs were found to peak 2.1 dB over the nominal normalized clock signal magnitude at a 30 GHz clock frequency. This peaking could potentially overdrive the clock buffers and registers, so resistors were added in series between the emitter followers of the clock drivers and the long microstrip transmission lines feeding the registers. These resistors absorb reflections on the microstrip transmission lines and reduce the clock overshoot to acceptable levels. As shown in Figure 2.4, adding series resistors in the clock distribution path reduces the overshoot down to acceptable levels. With the series resistors added, the best
case register input from the previous simulation now has the least clock bandwidth of all
the register inputs. For this example circuit, the targeted clock performance is 32 GHz,
so the addition of the series resistors does not negatively impact the circuit performance
because this register clock input still has 36 GHz of bandwidth. This method of clock
tree simulation was first reported in [51].

2.4 Emitter Coupled Logic

Historically, several forms of digital bipolar logic have been reported. These
include resistor-transistor logic (RTL), diode-transistor logic (DTL), transistor-transistor
logic (TTL), and emitter coupled logic. While TTL has been the most widely used,
particularly from the late 1960s to the late 1980s, ECL is the fastest [3]. ECL is based
on the bipolar differential pair with an emitter follower. An example of an ECL inverter
that uses the differential pair is shown in Figure 2.5. In ECL, the switching transistors
are kept in the active mode of operation, so the time needed to alter the base charge is
low, enabling high-speed operation [53] and current steering.

ECL uses a differential pair to steer the current (I_{EE}) depending on the input
differential signal (Ap/An). As described in Section 2.5 a sufficiently large input dif-
fferential, will result in almost all of the current steered through the transistor with the
highest input voltage. Thus, when the voltage at node Ap is sufficiently higher than An,
almost all of the current is steered through the left transistor in the differential pair, so
node Zp has a potential of approximately -I_{EE}*R_C volts. Since almost all of the cur-
rent is steered through the left side transistor, essentially no current flows though the
transistor on the right side of the differential pair, and node Zn is pulled up to the top
rail, which is 0 V in Figure 2.5. This convention is used for all of the circuits in this
thesis. The emitter followers on the outputs of the differential pair drive loads using the
differential signals Zp1 and Zn1. Without the emitter followers, differential pair alone
Figure 2.4: Example simulation of the clock distribution interconnects as microstrip transmission lines with line lengths, characteristic impedances, and loads estimated from the physical layout. The best case and worst case register inputs are shown both before and after the addition of series resistors to the clock distribution paths. Without the series resistors, there is some potential for overdrive in the worst case register inputs. The series resistors reduce the overdrive while maintaining a clock bandwidth well above the operating frequency of the circuit.
is only able to drive a low fan out of a couple of gates without losing performance and reducing the output differential across nodes Zp and Zn.

ECL can also be extended to multiple input logic gates. A two-input ECL gate is shown in Figure 2.6. In this two-input gate, two differential pairs are stacked upon each other. In each differential pair, the current is steered though the transistor with the highest voltage at the base node, Ap/An or Bp/Bn. The configuration shown in Figure 2.6 performs the AND operation. If the sense of the outputs is switched, so that nodes Zp and Zn are swapped, the logic gate in Figure 2.6 performs the NAND operation. From Demorgan’s theorem [54], when the sense of the both inputs and the output are swapped, the gate in Figure 2.6 functions as an OR gate. Likewise, if the sense of the both inputs are swapped, but the output is left as is, then the gate functions as a NOR gate. Thus, the two-input ECL gate topology shown in Figure 2.6 is very flexible. It can be used as an AND, NAND, OR, or NOR gate, just by modifying the sense of the inputs and outputs.
Other two-input configurations implemented in ECL, include the XOR/XNOR shown in Figure 2.7, and the latch shown in Figure 2.8. The current steering methodology can be extended to more than two inputs by stacking more differential pairs. For example, the three-input ECL AND gate shown in Figure 1.3.

### 2.4.1 Voltage Levels

In ECL, the logic levels are defined in terms of the number of diode drops below the top rail. Since the top rail is at ground in this thesis, the top voltage level goes from 0 to \(-I_{EE}*R_C\) volts. By convention, this range of voltages is defined as “level 0.” One diode drop (800 mV to 900 mV) below “level 0” is “level 1.” Another diode drop lower is “level 2.” Additional levels are possible, depending on the supply voltage and logic gates used.
Figure 2.7: ECL XOR logic gate schematic. The sense of the outputs can be swapped to achieve an XNOR gate from the topology.

Figure 2.8: ECL latch gate schematic. In this configuration, the latch is transparent when clkp has a higher voltage than clkn, and it is in latch mode when clkn has a higher voltage than clkp.
Generally, the uppermost transistors in the ECL gate transistor stacks are driven by “level 1” signals. Likewise, the next pair of transistors down in the stack are driven by “level 2” signals.

2.4.2 Current Mode Logic

It should be noted that in the literature, the term ECL is often used interchangeably with current mode logic (CML). In this thesis, the term ECL is used for logic gates with differential pairs followed by emitter followers. CML is used for the same gates, but without emitter followers. CML gates consume less power than ECL gates, since they eliminate the emitter follower circuitry. However, they do not have the ability to drive a large fan out or large node capacitances. Typically, CML performance will degrade after driving only two or three gates. The addition of emitter followers increases the drive capability of the logic gate and improves performance when there are multiple loads.

The logic levels that drive the CML gates are different from ECL, since the emitter followers are eliminated. The top differential pair is driven “level 0” instead of “level 1.” For multiple input gates, emitter followers are still needed to shift the logic level down the appropriate number of diode drops, so some of the power advantage of CML is negated.

In this thesis, the majority of the circuits are ECL. It should be noted that it is possible to drive the uppermost transistors of an ECL circuit with “level 0” signals. This is advantageous, particularly if a logic gate only drives a single load. In this case, the emitter followers on the gate can be eliminated to save power. This type of approach is not strictly ECL, and it is used only in limited circumstances. For the most part, however, the “standard” levels described in Section 2.4.1 are used.
2.5 Voltage Swing

For the ECL gates to operate properly, a sufficient voltage differential is needed at the inputs so that a majority of the current is steered through only one transistor of each differential pair. The voltage swing is the difference between the logic high and logic low signal, or $I_{EE} \times R_C$ volts in the case of the circuit in Figure 2.5. If the voltage swing is too low, the transistor in the differential pair that should be “off” will still have a significant amount of current. While the voltage swing can be made very large to ensure that the current is steered properly, the propagation delay of the logic gate will increase as the output voltage swing increases. Thus, the voltage swing must be balanced so that there is a sufficient differential to properly switch gates, but not an excessive differential that degrades performance.

The voltage swing design point is determined by simulating the current through the differential pair as a function of the input differential voltage. The test circuit for this simulation is shown in Figure 2.9. The transistor sizes and values of $R_C$ and $I_{EE}$ are typical of circuits used in the designs. This simulation uses a DC sweep of the differential voltage ($V_{diff}$), and the simulation is for circuit schematic only, with no extracted parasitic capacitors. The simulation output in Figure 2.10 shows the percentage of current through both transistors in the differential pair as a function of the input differential. From Figure 2.10, a voltage differential of 200 mV results in 99.25% of the current steered through the “on” transistor and a 300 mV differential results in 99.95% of the current steered through the “on” transistor. Below 200 mV, the percentage of current though the “on” transistor is rapidly reduced, and above 300 mV the increase in the percentage of current in the “on” transistor increases very slowly with $V_{diff}$. Thus, for a sufficient voltage swing that maximizes the gate performance, the internal voltage swing in the ECL circuits is kept within the 200 mV to 300 mV range.
Figure 2.9: Test circuit for designing voltage differential. The voltage differential is determined by $V_{\text{diff}}$, and $V_{\text{bias}}$ is used to bias the differential pair.

Figure 2.10: Simulation of the percentage of current through each leg of the differential pair as a function of voltage differential.
2.6 Bias Current

Providing a current source with a known, stable current is essential for circuit operation. Not only is the operating speed of the transistors dependent on the bias current, but the voltage swing is also dependent on the bias current. One simple way to generate a bias current is by using the simple current mirror shown in Figure 2.11. If the two transistors are equal, the input current \( I_{in} \) is mirrored to the output \( I_{out} \), such that

\[
I_{out} = \frac{I_{in}}{1 + \frac{2}{\beta}}. \tag{2.1}
\]

When \( \beta \) is large,

\[
I_{out} \approx I_{in}. \tag{2.2}
\]

In general, it is desirable to have the reference current on the right side of the circuit supply the bias voltage \( V_{bias} \) to several current sources. Since the transistor \( \beta \) is not infinite, the sum of all of the base currents will become significant, so that Equation 2.2 will not hold and \( I_{out} \) will not longer be equal to \( I_{in} \).

In this thesis, the current mirror is based on the beta helper design with emitter degeneration added. This current mirror is shown in Figure 2.12. On the left side
Figure 2.12: Current mirror with beta helper and emitter degeneration. The left side of the circuit generator the bias voltage. The bias voltage is used by multiple current sources in the circuit. An example current source is shown on the right side of the figure.

of Figure 2.12, the bias generator establishes the bias voltage used by current sources in the circuit. The right side of the figure shows an example of a current source that is used in the logic gates. The input current is established by using a resistive divider ($R_{bias}$). The additional transistor ($Q_b$) is the beta helper. In effect, $Q_b$ supplies the base current needed by the load transistors, allowing for multiple loads to be connected while maintaining an accurately mirrored current. The beta helper reduces the error in the current mirror by a factor of $(\beta + 1)$. When the transistors are identical,

$$I_{out} = I_{in} \left(1 - \frac{2}{\beta(\beta + 1)}\right).$$  \hspace{1cm} (2.3)

The current mirror also adds emitter degeneration in the form of resistors on the emitters of the transistors. Since there are process variations in implementing transistors, it is difficult to achieve perfect matching of the transistors. However, the emitter degeneration significantly improves current matching [55] by compensating for process mismatches.
In this thesis, the current mirror based on Figure 2.12 is used as a bias generator in all of the designs except for the resistor-only current source accumulator that is described in Section 3.4.2. The resistor values are modified and a different number of diodes may be used depending on the supply voltage of a particular design. In individual circuits, when a current source is shown, it is actually implemented as a transistor with its base connected to the $V_{\text{bias}}$ node of a current source and with an emitter degeneration resistor, as shown on the right side of Figure 2.12. Individual current sources can be modified by changing the size of the transistor relative to the mirror transistor in the bias generator. The degeneration resistor is also scaled to maintain a constant degeneration voltage.

2.7 Current Source Output Resistance

Ideally, the current source outputs a constant current that is independent of the voltage across its collector and emitter. However, these current sources are not ideal. In the extreme case, the collector and emitter voltages may be so close that the output current is driven to 0 mA. At higher collector voltages, the current becomes linearly dependent on the collector voltage. This linear dependence is expressed as the output resistance.

Using the bias generator and a current source that both have a supply voltage of -3.8 V, the collector voltage of the current source transistor is swept using a DC simulation in Spectre. The output of this simulation shows the dependence of the output current on the collector voltage. It is shown in Figure 2.13. The DC sweep goes up to only -1.8 V, because collector voltages higher than this are avoided to ensure that the transistors are in the safe operating area.

Below -3.2 V, as shown in Figure 2.13, the bias transistor is not fully on, and the current is very highly dependent on the collector voltage. In the circuits in this thesis, the current sources are not operated in this region. In fact, they are generally
Figure 2.13: DC sweep simulation of the current source current versus the collector voltage. The output resistance is 3.7 k\(\Omega\) over the linear region.
kept above -2.8 V. At -2.8 V and above, the relationship between the current and voltage is approximately linear. The slope of the current line in this region gives the output resistance ($R_{out}$), such that

$$R_{out} = \frac{\Delta V_{\text{collector}}}{I_{\text{bias}}}.$$  

(2.4)

For this circuit, $R_{out}$ is 3.7 kΩ. In this thesis, not all of the circuits will have the same collector voltage at the current source, so it is important to factor in the output resistance when biasing circuits.

### 2.8 Conclusion

This chapter provided an overview of some of the general aspects of circuit design in this thesis. Many of these aspects are common to multiple circuit designs in this thesis. Accumulator and DDS circuits will be presented in the following chapters.
CHAPTER 3

Design of Adders and Accumulators in InP

In this chapter, several accumulator designs will be discussed. All of the accumulators are designed so that they can potentially be integrated as phase accumulators in DDS circuits. While the accumulators discussed in Section 3.2.3 through Section 3.2.6 are not integrated into DDS circuits, they are fabricated inside test circuits containing a DAC. The DAC is included to simplify testing by combining several high-speed digital accumulator outputs into a single high-speed analog output. The accumulators discussed in Section 3.3.1 and Section 3.3.2 are integrated and fabricated as phase accumulators in DDS test circuits. Since these accumulators drive phase to sine converters, no direct measured results of these accumulators are available. However, the operating speed of these accumulators is inferred from the DDS circuit measurements. These DDS circuits will be discussed in further detail in Chapter 4.

Two additional accumulator designs are discussed in Section 3.4. The first design uses a triple-tail circuit approach is discussed in Section 3.4.1. This approach is not implemented as a complete accumulator circuit even though others have reported advantages using a triple-tail approach for high-speed lower power designs [56, 57], because the simulation results do not show an improvement over the other designs. The second design is described in Section 3.4.2. Instead of using the current mirror as a current source, it uses resistor-only current sources. This approach saves a significant amount of power, but may be risky. At the time of the thesis publication, it is still in fabrication.

Before the accumulator designs can be discussed in detail, it is necessary to review some of the basics of adders and accumulators. Much of the information in the following section can be found in an introductory digital logic text, such as R. Katz [54], or a more advanced text, such as D. Hodges et al. [11].
Table 3.1: Truth table for the full adder building block. A and B are the two input bits, \(C_{in}\) is the carry input, S is the sum, and \(C_{out}\) is the carry output.

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3.1 Review of Adder and Accumulator Basics

This section discusses some of the basics of adders and accumulators. The full adder is introduced as the primary building block for multiple bit adders. Basic multiple bit adder architectures including the carry ripple adder, carry look ahead adder, carry select adder, and pipelined adders are then discussed. Finally, an overview of the unique attributes of accumulator circuits in comparison to adder circuits is discussed.

3.1.1 Full Adder

The full adder is the basic building block in many adder designs. It performs the operation of adding two bits (A and B), combined with a carry input (\(C_{in}\)). It outputs both the sum (S) and carry output (\(C_{out}\)) of the addition operation. The truth table for the full adder is shown in Table [3.1].

As shown in Table [3.1], the output of the full adder is the 2-bit sum of three 1-bit inputs, where the \(C_{out}\) is the most significant bit (MSB) and S is the least significant bit (LSB) of the 2-bit sum. The carry output is high whenever two or three of the inputs (a majority of the inputs) are high. Logically, this is

\[
C_{out} = A \cdot B + A \cdot C_{in} + B \cdot C_{in} = A \cdot B + C_{in} \cdot (A + B). \tag{3.1}
\]
The sum output is high when either one or three of the inputs are high. Logically, the this is

\[ S = A \oplus B \oplus C_{in}. \]  \hspace{1cm} (3.2)

### 3.1.2 Carry Ripple Adder

The simplest type of multiple-bit adder is formed by stringing together multiple full adders. An example of a 4-bit carry ripple adder is shown in Figure 3.1. The clock signal latches the sum outputs at the end of the addition operation. The carry outputs occur within a single clock cycle. Although a 4-bit example is shown, the concept can be extended to larger bit-widths. This circuit is called the carry ripple adder, because changes in the carry bits “ripple” or propagate from the LSB to the MSB of the adder.

The carry ripple adder is slow, especially for large bit-widths. The worst case propagation delay of the carry ripple adder is dependent on the propagation delay of the carry circuit \( t_{carry} \), the propagation delay of the sum \( t_{sum} \), and the bit width. In general, it is given by the propagation delay through all but the MSB carry, plus the MSB sum propagation delay, or

\[ t_{carry\text{ripple}} = (n - 1) \cdot t_{carry} + t_{sum}. \]  \hspace{1cm} (3.3)
For a 4-bit carry ripple adder, the propagation delay is \(3 \cdot t_{\text{carry}} + t_{\text{sum}}\). While many circuit implementations are possible, for simplicity it will be assumed that only two-input logic gates are available and that all two-input gates have a propagation delay of \(t_{\text{gate}}\). In reality, and as will be shown later in this thesis, this assumption does not always hold. The gate delay is dependent on the technology and design approach used, so it is not necessarily equal for the sum and carry circuits. Also, it is possible to use three-input gates which may have a different propagation delay than two-input gates. Despite the imperfections in the assumption, it is useful for making a rough comparison to other adder architectures.

Continuing with the rough assumption of a \(t_{\text{gate}}\) propagation delay for all two-input gates, the carry \((t_{\text{carry}})\) and sum \((t_{\text{sum}})\) propagations delays are both \(2 \cdot t_{\text{gate}}\), since these each of these gates are comprised of a pair of two-input gates. As a result, the 4-bit carry ripple adder has a worst case propagation delay of \(8 \cdot t_{\text{gate}}\). This worst case propagation delay increases linearly with bit-width.

### 3.1.3 Carry Lookahead Adder

An alternative to the carry ripple adder that reduces the propagation delay is the carry lookahead adder. To reduce dependence on the carry propagation delay, the full adders in the carry lookahead adder produce “generate” and “propagate” signals independently of the \(C_{\text{in}}\) input. The generate signal \((G)\) is high whenever the full adder block would have a carry out, regardless of the state of \(C_{\text{in}}\). The generate signal is

\[
G = A \cdot B. \tag{3.4}
\]

The propagate signal \((P)\) is high if and only if the full adder block would generate a carry output when \(C_{\text{in}}\) is high. The propagate signal is

\[
P = A \oplus B. \tag{3.5}
\]
For the $i^{th}$ stage of multi-bit wide adder, the carry input is $C_i$ and the carry output is $C_{i+1}$. The sum and carry output at the $i^{th}$ stage of an adder are

$$S_i = P_i \oplus C_i,$$

$$C_{i+1} = G_i + P_i \cdot C_i.$$  

(3.6)  

(3.7)

Assuming that the carry input to a 4-bit adder is a logic low ($C_0 = 0$), the propagate and generate signals can be used to compute the sum such that:

$$S_0 = P_0$$

(3.8)

$$S_1 = P_1 \oplus G_0$$

(3.9)

$$S_2 = P_2 \oplus (G_1 + P_1 \cdot G_0)$$

(3.10)

$$S_3 = P_3 \oplus (G_2 + P_2 \cdot G_1 + P_2 \cdot P_1 \cdot G_0).$$

(3.11)

The carry lookahead adder is not dependent on the carry rippling through every full adder circuit, so its propagation delay is dependent on the propagation delay of the most complex logic function in the circuit. An example of a 4-bit carry lookahead adder is shown in Figure 3.2.

In the 4-bit adder example, the propagation delay is dependent on the logic used to implement $S_3$. Using the rough assumption for comparison from Section 3.1.2 that each two-input gate has a propagation delay of $t_{gate}$, the 4-bit carry lookahead adder has a worst case propagation delay of $5 \cdot t_{gate}$. For larger bit-widths, the carry lookahead adder can be broken up into groups of 4-bit carry lookahead adders that each have a group propagate and group generate signal. For a 16-bit carry lookahead adder, the worst case propagation delay is $10 \cdot t_{gate}$. This is much faster than a 16-bit carry ripple adder, which has a worst case propagation delay of $32 \cdot t_{gate}$.  

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3.1.4 Pipelined Adders

Since the propagation delay determines the maximum clock frequency at which the adder can be operated, a reduced propagation delay will lead to a faster adder. The propagation delay cannot always be reduced in a technology, but it can be reduced in effect by using pipelining. In pipelining, latches are inserted in between sections of logic within the adder. This effectively reduces the total adder propagation delay to the propagation delay between latches and allows for the clock speed to be increased.

Pipelining can be integrated into any adder architecture by inserting clocked pipeline registers in between sections of logic. Thus, a carry ripple adder or a carry lookahead adder can be implemented with pipelining. A potential disadvantage of using pipelining in an adder is that pipelining introduces latency between the input of data and when the specific result dependent on the data is output. The latency is determined by the number of stages, or slices of logic between latches, in the adder. Depending on the application, latency may or may not be a problem. In pipelined DDS applications, for example, an increase in latency decreases DDS agility. However, this is usually not a problem because the typical decreases is agility are generally acceptable for most applications.
Despite the latency inherent in the pipeline approach, the maximum clock frequency is increased in comparison to other adders without pipelining. This is because a sum is output every clock cycle, while partially computed sums are stored within the pipeline latches. In the extreme case, latches can be inserted between every logic gate so that the clock speed is dependent on the propagation delay of a single logic gate plus the propagation delay of the latch. As the number of pipeline stages increases, so does the power consumption of the adder. Not only is extra power needed for the latches, but it is also required by the circuitry that supplies a clock signal to the latches.

3.1.5 Accumulators

Accumulators are a special case of adders. They are typically formed by feeding the latched sum output of the adder back into the B input. For convenience, the discussion of accumulation feedback will deal with the feedback into the B input, but it is possible for the A input to be used instead. During each clock cycle, the adder sums the accumulation increment (A) and the previous sum (B) to operate as an accumulator.

In a DDS, the A input typically changes at a much lower frequency than the clock, because it is used as the frequency control word (FCW). Thus, while the circuitry associated with the B input must be capable of operating at the clock frequency, the circuitry associated with the A input can have lower performance. The DDS phase accumulator can be configured so that some of the performance on the A input circuitry is traded for improved performance on the B input circuitry. Since the accumulator performance is limited by the B input circuitry, the accumulator operating frequency is increased. This allows for a higher accumulator maximum frequency compared to an equivalent adder design.

Power can also be reduced in an adder configured as an accumulator. In a pipelined adder, the data must be buffered both before and after the accumulation operation is completed to ensure that the pipelined data is in the proper clock cycle of
buffering. As described in [58], the accumulator can feed the sum output back locally to the adder blocks, eliminating the need for the pre-buffering registers on the A inputs. This modification cuts the number of buffering registers in half, greatly reducing the power.

In the following sections, several accumulator designs will be discussed in depth.

### 3.2 High-Speed 4-bit Accumulators

The first set of accumulators are designed with high-speed operation as the main design goal. All of the accumulators have a bit-width of 4-bits. This short bit-width allows for demonstrating the operation of the accumulators while keeping a relatively low transistor count. This helps to minimize problems with power, heat, clock distribution, and reliability, so that the basic circuit designs are the main focus.

In previous work in InP by T. Mathew et al. [27, 59], only accumulator components were implemented, not a complete accumulator circuit. The previous work used four-level series-gated logic merged with latches to implement carry and sum circuits. This design methodology is described further in Section 3.2.3. The previous work demonstrated a carry test circuit that was configured as a divide by two circuit operating at a maximum clock frequency of 19 GHz.

Four accumulators are presented in this section. In Section 3.2.3, ACCV1 is a 4-bit accumulator based on the design reported by T. Mathew et al. [27, 59]. Like the previous work, a carry test circuit is implemented as a divide by two circuit for comparison. This accumulator also extends upon the previous work by implementing a full 4-bit accumulator, instead of just the accumulator components.

The four-level series-gated sum and carry circuits in ACCV1 require a large supply voltage, but many of the other circuits in the accumulator do not. Particularly in the latch circuits, which will become a large portion of the circuit as bit-widths are extended, the large supply voltage results in wasting power. In Section 3.2.4, ACCV2 is
presented as a step towards reduced power consumption. This accumulator uses a single-level parallel-gated carry circuit instead of the four-level series-gated carry circuit used in ACCV1. Since the sum circuit is not changed, the supply voltage and power are not reduced compared to ACCV1. However, the ACCV2 design is useful as a test bench for the new carry circuit and for a comparison to the previous design, since only the carry circuit is changed.

Reduced power accumulator designs ACCV3 and ACCV4 are presented in Section 3.2.5 and Section 3.2.6. These accumulators use a three-level series-gated sum circuit that allows for the reduction of the supply voltage by a diode drop. This power supply reduction leads to a reduced power consumption. ACCV3 and ACCV4 have a similar architecture, but some layout and clock buffering differences. ACCV3 uses a layout and clock buffering that is similar to the ACCV1 and ACCV2 designs, so it provides a more direct comparison. ACCV4 uses a more compact layout topology and stronger clock buffers, so that a more compact layout can be investigated and a more conservative approach to clock buffering is used.

### 3.2.1 Test Circuits

In the previous work reported by T. Mathew et al. [27, 59], the carry circuit was the limiting factor for accumulator performance because it was slower than the sum circuit. To estimate an upper bound for accumulator performance, the performance limiting carry circuit is tested as a divide by two circuit. The divide by two test circuit also provides a test of the four-level series-gated carry circuit functionality.

The divide by two carry test circuit is shown in Figure 3.3. It is used for testing the four-level series-gated and single-level parallel-gated carry circuits in ACCV1 and ACCV2 in Section 3.2.3 and Section 3.2.4. It uses a carry circuit with a logic high on the ‘A’ input, a logic low on the ‘B’ input, and with the carry output inverted and fed back into the ‘C’ input. Since the ‘A’ input is high and the ‘B’ input is low, the carry
Figure 3.3: Block diagram of divide by two test circuit.

Figure 3.4: Block diagram of the 4-bit accumulator test circuit. The DAC combines the four high-speed digital sum outputs of the accumulator (S(3:0)) into a single high-speed analog output. This output can be observed on a sampling oscilloscope.

circuit tracks the ‘C’ input. The inversion of the carry output into the ‘C’ input results in a divide by two circuit. A buffer drives the divide by two signal off of the test chip.

In addition to the divide by two test circuits for ACCV1 and ACCV2, all four of the 4-bit accumulator test circuits contain a 4-bit DAC for testing purposes. An example of a 4-bit accumulator test circuit is shown in Figure 3.4. The design allows for any 4-bit increment (0 through 15) to be input into the accumulator in A(3:0). Since the accumulator has four high-frequency outputs (S(3:0)), the on-chip 4-bit DAC is used to generate a single high-speed analog output. The DAC output preserves all of the information for determining proper operation in an analog output that can be observed on a sampling oscilloscope.
3.2.2 Measurements

The divide by two carry test circuits and accumulator test circuits are tested on-wafer using an Alessi probe station with HP and Agilent test equipment. DC inputs are brought onto the device under test (DUT) with multi-pin probe cards, and RF inputs and outputs are interfaced with high-frequency ground-signal-ground-signal-ground (GS-GSG) probes. An additional frequency multiplier is needed for input clock frequencies above 50 GHz for the carry divide by two test circuits. The test setup for clock frequencies below 50 GHz is shown in Figure 3.5, and the test setup for test frequencies above 50 GHz is shown in Figure 3.6.

3.2.3 Accumulator ACCV1

A design based on the work of T. Mathew et al. [27, 59] is implemented as a baseline 4-bit accumulator, named ACCV1, in the InP DHBT technology described in
Section 2.1. Unlike traditional pipelined designs, the ACCV1 design merges combinational logic structures with latches, using four-level series-gated logic. The four-level series-gated carry and sum circuits are shown in Figure 3.7 and Figure 3.8, respectively. The four-level series-gated logic leads to high power consumption compared to an approach with separate logic and latches, because it is necessary to add an extra diode drop to the power supply in order to support the clock transistors added to the stack of transistors. However, overall circuit performance is increased in comparison to approaches with separate logic and latches because the front-end buffers of the latches are replaced with carry and sum logic, reducing the overall propagation delay.

Unlike the T. Mathew et al. design [27, 59], which only implements separate sum and carry circuits, the ACCV1 is a full 4-bit accumulator. The 4-bit accumulator is pipelined, and it is built from 2-bit adder blocks and 2-bit register blocks. The ACCV1 4-bit accumulator is shown in Figure 3.9. The 2-bit adder blocks are formed by combining
carry circuits, sum circuits, and latches, as shown in Figure 3.10. The components labelled ‘carry & latch’ and ‘sum & latch’ refer to the circuits in Figure 3.7 and Figure 3.8 that merge the logic with latches. The pipelined structure of the 4-bit accumulator in Figure 3.9 can be easily expanded to an arbitrary 2N-bit width [58]. This is particularly useful for DDS designs, where larger bit widths are typically required for improved SFDR and finer frequency resolution. Since the accumulator is pipelined, an expanded data word size can be achieved with little impact on high frequency performance.

### 3.2.3.1 Simulation Results

The divide by two carry test circuit simulates up to a maximum clock frequency of 55 GHz. An output of the divide by two carry test circuit simulation is shown in Figure 3.11. In simulation, the 4-bit accumulator ACCV1 operates up to a maximum clock rate of 43 GHz. In the ACCV1 simulation results shown in Figure 3.12, only the outputs of the accumulator prior to the DAC are shown in order to more clearly
Figure 3.9: Block diagram of the pipelined 4-bit accumulator using 2-bit adders and 2-bit registers.

Figure 3.10: Block diagram of the 2-bit adder comprised of carry, sum, and latch circuits.
<table>
<thead>
<tr>
<th>Circuit</th>
<th>Sub-Circuits</th>
<th>Power</th>
</tr>
</thead>
<tbody>
<tr>
<td>Accumulator Core</td>
<td></td>
<td>3.008 W</td>
</tr>
<tr>
<td></td>
<td>2-Bit Accumulator</td>
<td>829 mW</td>
</tr>
<tr>
<td></td>
<td>2-Bit Register</td>
<td>550 mW</td>
</tr>
<tr>
<td></td>
<td>Clock Circuitry</td>
<td>800 mW</td>
</tr>
<tr>
<td>Support Circuitry</td>
<td>Input Buffers</td>
<td>1.163 W</td>
</tr>
<tr>
<td></td>
<td>Output Stage/DAC</td>
<td>472 mW</td>
</tr>
<tr>
<td></td>
<td>Bias Generators</td>
<td>600 mW</td>
</tr>
<tr>
<td></td>
<td></td>
<td>91 mW</td>
</tr>
<tr>
<td>Total</td>
<td></td>
<td>4.171 W</td>
</tr>
</tbody>
</table>

Table 3.2: Simulated power consumption breakdown for the ACCV1 accumulator test circuit.

illustrate the proper operation of the 4-bit accumulator. It should also be noted that the test chips have differential internal circuitry, but only single-ended outputs are shown in Figure 3.12 for clarity.

The accumulator core, which consists of two 2-bit accumulators, a 2-bit register, and the associated clock circuitry consumes 3.008 W of power in simulation. Other circuitry needed for the accumulator test circuit, such as the DAC and input buffers uses 1.163 W. The ACCV1 accumulator test circuit uses a total of 4.171 W. A breakdown of the simulated power consumption is shown in Table 3.2.

### 3.2.3.2 Measurement Results

The divide by two carry test circuit and the accumulator test circuit were fabricated in both TC6 and TC7 design runs. They were tested using the test setup described in Section 3.2.2. The test results are described below.

#### 3.2.3.2.1 TC6 Measurement Results

The carry test circuit was found to operate up to a maximum clock frequency of 52 GHz. The measured output signal of 26 GHz, generated from a 52 GHz clock is shown in Figure 3.13. This output was captured using a high-speed sampling oscilloscope. There is some attenuation in the signal compared to
Figure 3.11: Simulation of the four-level parallel-gated carry test circuit. The test circuit uses the carry circuit as a divide by two circuit to estimate an upper bound for the accumulator operating frequency. This simulation includes extracted parasitic capacitors and is shown at the maximum operational simulated clock frequency of 55 GHz.
Figure 3.12: Simulation of the ACCV1 4-bit accumulator using four-level series-gated logic. The plot shows the outputs of the four sum bits. The simulation includes extracted parasitic capacitors and is shown at the maximum operation frequency of 46 GHz.
the simulation result because of extra parasitic capacitance that is added by the probes and cables that connect the test chip to the test and measurement equipment. This result is only 3 GHz, or 5.5% lower than the maximum simulated clock frequency. This difference could be due to process variations or inaccuracies in the models. The VIP-2 process was being modified during fabrication run TC6, so the model accuracy is questionable.

It should be noted that circuit yield in this particular fabrication run was low, so there were very few working test sites on-wafer for comparison to simulation data. Shorts in the capacitors were a particular problem. A table of the yield for the divide by two test circuit is shown in Table 3.3. Low yield in this fabrication run also resulted in no working ACCV1 4-bit accumulator test circuits. Failure modes included shorted capacitors, circuits that were under current, circuits with correct currents, but no outputs, and circuits with incorrect outputs. A table of the yield for the ACCV1 accumulator test circuit is shown in Table 3.4. Since the working carry test circuits were close to the simulation frequency, it is expected that ACCV1 would have performed near the 43 GHz clock rate expected from simulation had the yield been higher.
Figure 3.13: High-speed sampling oscilloscope screen capture of the 26 GHz output from the four-level series-gated carry test circuit operating as a divide by two circuit with a 52 GHz clock. This output is attenuated compared to the simulation results because of additional parasitic capacitance from the probes and cables that connect the test chip to the test equipment.
Table 3.5: Yield for the four-level series-gated carry divide by two and ACCV1 accumulator test circuits in fabrication run TC7. A device is considered to pass if it operates correctly above a 24 GHz clock rate.

<table>
<thead>
<tr>
<th>Result</th>
<th>Divide by Two</th>
<th>ACCV1</th>
</tr>
</thead>
<tbody>
<tr>
<td>Fail (Short)</td>
<td>1</td>
<td>6</td>
</tr>
<tr>
<td>Fail (Other)</td>
<td>3</td>
<td>8</td>
</tr>
<tr>
<td>Pass</td>
<td>6</td>
<td>3</td>
</tr>
</tbody>
</table>

3.2.3.2.2 TC7 Measurement Results  The four-level series-gated carry divide by two test circuit and the ACCV1 test circuit were also fabricated as part of fabrication run TC7. The were tested in the same manner used for the TC6 fabrication run. The divide by two carry circuit operated up to a maximum clock frequency 51 GHz, which is comparable to the 52 GHz maximum measured from TC6. This is still slower than expected from simulation. Three ACCV1 devices were operational in the TC7 fabrication run. The fastest ACCV1 test circuit operated up to a maximum clock frequency of 38 GHz and was measured to consume 3.8 W of power. This is slower than expected from simulation, however, other test circuits were also found to operate slower than expected on this fabrication run. Output waveforms were not captured for either test circuit. The yield for both the divide by two and ACCV1 test circuits is shown in Table 3.5. Capacitor shorts were again a problem in this fabrication run, particularly in the ACCV1 test circuit, which had a greater number and larger capacitors than the divide by two test circuit. A device is considered to pass if it operates correctly above a 24 GHz clock rate.

The measured power for ACCV1 was 3.8 W, including all of the circuitry such as the DAC. This is lower than the simulated power consumption of 4.171 W. Discrepancies between the models and the fabricated devices could explain this difference. The internal components of ACCV1 do not have separate power supplies, so the power breakdown must be estimated. Using the ratios established Table 3.2 these estimates
can be determined. The total consumption of the 4-bit accumulator is 2.76 W. The 4-bit accumulator has two 2-bit adders consuming 0.76 W of power each, a 2-bit register consuming 0.51 W of power, and clock tree circuitry that consumes 0.73 W of power. If the clock tree is partitioned so that the adders and registers include the power from the clock tree circuitry that they require, the 2-bit adders would use 1.00 W and the 2-bit registers would use 0.76 W.

3.2.4 Accumulator ACCV2

The use of four-level series-gated logic in ACCV1 described in Section 3.2.3 allows for high-speed operation, but it is achieved at the expense of high power consumption. The power supply voltage must be high enough to support the output voltage swing, three diode drops from the logic, one diode drop from the clock, a diode drop from the current source, and the voltage drop across the current source degeneration resistor. In this technology, this requires a -5.5 V supply voltage. In some circuitry only two total diode drops are actually required for operation. As a result, power is wasted when these circuits are connected to a high supply voltage. If even one diode drop could be removed from the power supply overhead, then substantial power savings can be realized. The latch and register circuits are an example of a circuit that requires only two diode drops. In high-speed pipeline accumulators, the latch and register circuits comprise a significant portion of the total circuitry. To move towards an accumulator with a lower supply voltage, the four-level series-gate carry circuit shown in Figure 3.7 is replaced with the single-level series-gated carry circuit [50] shown in Figure 3.14. A patent application was filed for the single-level parallel-gated carry circuit [60].

Single-level parallel-gated logic is well suited for the carry circuit because it performs a majority operation, essentially detecting when two or three of the inputs are high. When all three inputs are either high or low, a full differential is seen across Xp and Xn, since all of the current is steered through one leg of the circuit. When one or
two of the inputs are high, the differential across Xp and Xn is reduced, since 1/3 of the current is steered through one leg of the circuit while 2/3 of the current is steered through the other leg of the circuit. Although this method has a reduced differential across Xp and Xn for some input states, this signal is sampled by the latch and a full differential is generated for propagation to subsequent circuitry. Figure 3.15 shows a simulation of the single-level parallel-gated carry circuit schematic illustrating the output voltage at Xp and Xn of the carry for possible input combinations, as well as the full differential output after the first stage of the latch.

Like ACCV1 described in Section 3.2.3, ACCV2 uses the four-level series-gated sum circuit shown in Figure 3.8. As a result, the supply voltage can not be reduced in ACCV2. In this implementation, the single-level parallel-gated carry circuit uses more current than the four-level series-gated carry circuit, so overall power consumption in ACCV2 is actually slightly higher than the previous design. However, ACCV2 provides an intermediate design point where only the carry circuit is changed compared to
Figure 3.15: Simulation of the output of the single-level parallel-gated carry circuit. The upper plot shows the areas with reduced differential for the $X_p$ and $X_n$ outputs for the states where 0, 1, 2, and 3 inputs are high. The lower plot illustrates how the full differential is regained by buffering from the latch.
ACCV1. The single-level parallel-gated carry circuit illustrates a proof of concept for reducing power in later designs. If the sum circuit that is implemented in ACCV2 is replaced by a circuit using fewer levels, then the accumulator can have a lower supply voltage and operate at a lower power. The power reduction associated with the removal of one diode drop from the supply voltage is expected to be around 15%. Accumulators with alternative sum circuit designs that allow for lower power supply voltages and reduced power consumption are discussed in later sections.

3.2.4.1 Simulation Results

As in Section 3.2.3, a carry test circuit and an accumulator test circuit are both simulated. The single-level parallel-gated carry test circuit simulates as a divide by two circuit at a maximum clock frequency of 52 GHz. This simulation output is shown in Figure 3.16. The result is slower than the simulation from the four-level series-gated design by only 3 GHz, or 5.5%.

The ACCV2 4-bit accumulator test circuit simulates up to a maximum clock frequency of 46 GHz. A plot of the four sum output bits from the accumulator at the maximum frequency of 46 GHz is shown in Figure 3.17. Overall, the accumulator simulated 3 GHz faster than the previous design. This is unexpected, since the carry test circuit is slower and the carry should be the limiting factor on high-speed operation. Even though ACCV1 and ACCV2 were very similar designs, subtle layout differences between the two carry designs may have potentially led to enough variation in parasitic capacitance to account for this difference.

The accumulator core, which consists of two 2-bit accumulators, a 2-bit register, and the associated clock circuitry uses 3.320 W in simulation. Other circuitry needed for the accumulator test circuit, such as the DAC and input buffers uses 1.163 W. The ACCV2 accumulator test circuit uses a total of 4.483 W. This is higher than the 4.171 W
Figure 3.16: Simulation of the single-level parallel-gated carry test circuit. The test circuit uses the carry circuit as a divide by two circuit to provide an estimate of the upper bound of the accumulator operating frequency. This simulation includes parasitic extracted capacitors and is shown at the maximum operation frequency of 52 GHz. The simulation shows the output at the output pad of the extracted test chip layout, so the extracted parasitic capacitances affect the waveform shape.
Figure 3.17: Simulation of ACCV2 showing the outputs of the four sum bits. The simulation includes parasitic extracted capacitors and is shown at the maximum operation frequency of 46 GHz.
simulated power consumption of ACCV1, since the single-level parallel-gated carry circuit uses more current sources and the same supply voltage. However, as stated earlier, ACCV2 provides a proof of concept for the single-level parallel-gated carry circuit, and it is not intended to be a lower power stand-alone accumulator. A breakdown of the simulated power consumption for the ACCV2 accumulator is shown in Table 3.6.

<table>
<thead>
<tr>
<th>Circuit</th>
<th>Sub-Circuits</th>
<th>Power</th>
</tr>
</thead>
<tbody>
<tr>
<td>Accumulator Core</td>
<td></td>
<td>3.320 W</td>
</tr>
<tr>
<td>2-Bit Accumulator</td>
<td>985 mW</td>
<td></td>
</tr>
<tr>
<td>2-Bit Register</td>
<td>550 mW</td>
<td></td>
</tr>
<tr>
<td>Clock Circuitry</td>
<td>800 mW</td>
<td></td>
</tr>
<tr>
<td>Support Circuitry</td>
<td></td>
<td>1.163 W</td>
</tr>
<tr>
<td>Input Buffers</td>
<td>472 mW</td>
<td></td>
</tr>
<tr>
<td>Output Stage/DAC</td>
<td>600 mW</td>
<td></td>
</tr>
<tr>
<td>Bias Generators</td>
<td>91 mW</td>
<td></td>
</tr>
<tr>
<td>Total</td>
<td></td>
<td>4.483 W</td>
</tr>
</tbody>
</table>

Table 3.6: Simulated power consumption breakdown for the ACCV2 accumulator test circuit.

### 3.2.4.2 Measurement Results

The circuits were fabricated as part of the TC6 fabrication run. Both the carry and accumulator test circuits were tested using the same test setup as the ACCV1 test circuits. These test setups are described in further detail in Section 3.2.2.

A microphotograph of the ACCV2 divide by two carry test chip is shown in Figure 3.18. The carry test circuit was measured to operate up to a maximum output frequency of 27.5 GHz with a clock frequency of 55 GHz, and is shown in Figure 3.19. This is faster than the simulation frequency of 52 GHz by 3 GHz. Yield for the carry test single-level parallel-gated circuit is shown in Table 3.7.

A microphotograph of the ACCV2 4-bit accumulator test chip is shown in Figure 3.20. The sum bits from the accumulator could not be directly measured, however, the output of the DAC demonstrates proper operation of the accumulator. Two examples
Figure 3.18: Microphotograph of the ACCV2 single-level parallel-gated carry test chip. The chip is 1220 µm by 1025 µm.

Figure 3.19: Oscilloscope screen capture of the carry test circuit output at 27.5 GHz with a 55 GHz clock frequency.
Table 3.7: Yield for the four-level series-gated carry divide by two test circuit in fabrication run TC6.

<table>
<thead>
<tr>
<th>Result</th>
<th>Number of Devices</th>
</tr>
</thead>
<tbody>
<tr>
<td>Fail (Short)</td>
<td>4</td>
</tr>
<tr>
<td>Fail (Correct Power, No Output at 50 GHz)</td>
<td>3</td>
</tr>
<tr>
<td>Pass (Operational for 50 GHz Clock or Above)</td>
<td>5</td>
</tr>
</tbody>
</table>

Table 3.8: Yield for the ACCV2 accumulator test circuit in fabrication run TC6.

<table>
<thead>
<tr>
<th>Result</th>
<th>Number of Devices</th>
</tr>
</thead>
<tbody>
<tr>
<td>Fail (Short)</td>
<td>8</td>
</tr>
<tr>
<td>Fail (Low Current)</td>
<td>3</td>
</tr>
<tr>
<td>Fail (Correct Current, No Output)</td>
<td>1</td>
</tr>
<tr>
<td>Fail (Incorrect Output)</td>
<td>1</td>
</tr>
<tr>
<td>Pass</td>
<td>1</td>
</tr>
</tbody>
</table>

of measured accumulator output sequences are shown in Figure [3.21] and Figure [3.22]. In the first example, an input increment of 7 is used to create the digital sequence of 15, 6, 13, 4, . . . In this configuration, the 16 discrete DAC output voltage levels are clearly illustrated in Figure [3.21], while operating with a clock frequency of 41 GHz. In the second example, the accumulator is configured as a divide by two circuit by using an input increment of 8. In this case, the most significant bit of the 4-bit digital value changes every clock cycle, creating a 20.5 GHz output signal from a 41 GHz clock frequency, as shown in Figure [3.22]. The 41 GHz maximum operation frequency was lower than the expected operating frequency of 46 GHz. However, yield was low for the ACCV2 test circuit in this fabrication run (TC6), with only one operational test site out of 14 tested sites. The yield for the ACCV2 test circuit is shown in Table [3.8]. Therefore, there is not enough information to determine if the performance difference was due to circuit design or if the single operational accumulator was on a “slow” test site.

The measured power for ACCV2 was 4.1 W, including all of the circuitry such as the DAC. This is lower than the simulated power consumption of 4.483 W. Discrepancies between the models and the fabricated devices could explain this difference.
Figure 3.20: Microphotograph of the ACCV2 test circuit. The chip is 1725 µm by 1025 µm.

Figure 3.21: Oscilloscope screen capture of the DAC output of ACCV2 4-bit accumulator test circuit with 41 GHz clock frequency and input increment of 7. The digital output sequence is labelled on waveform.
Figure 3.22: Oscilloscope screen capture of the DAC output of the ACCV2 4-bit accumulator test circuit with 41 GHz clock frequency and input increment of 8 acting as a divide by two circuit with a 20.5 GHz output.

The internal components of ACCV2 do not have separate power supplies, so the power breakdown must be estimated. Using the ratios established Table 3.6, these estimates can be determined. The 4-bit accumulator is estimated to consume 3.04 W, with two 2-bit adders at 0.90 W each, a 2-bit register at 0.51 W, and clock tree circuitry at 0.73 W. If the clock tree is divided so that the adders and registers include the power from the clock tree circuitry that they require, the 2-bit adders would use 1.14 W and the 2-bit registers would use 0.76 W. The 41 GHz ACCV2 design is reported in [50].

3.2.5 Accumulator ACCV3

The single-level parallel-gated carry circuit in Section 3.2.4 takes a step towards reduced power consumption since it can use a lower power supply voltage than is required for designs using four-level series-gated logic. However, a lower supply voltage can not be used for the ACCV2 circuit since the supply voltage is still constrained by the four-level series-gated sum circuit. Alternative sum circuits with fewer series gates
would allow for a power supply voltage reduction when used in conjunction with the single-level parallel-gated carry circuit. One such alternative is the three-level series-gated sum circuit [61] shown in Figure 3.23 and incorporated in accumulator ACCV3.

While this sum circuit topology still has logic merged with a latch, it uses an additional stage of logic, so that one of the diode drops can be eliminated. In general, the addition of an extra stage of logic leads to a longer propagation delay and a lower maximum operating frequency. However, in the accumulator the relation of the inputs to the sum circuit allows for operation at roughly the same speed as the four-level series-gated sum circuit. In the first stage of the sum circuit, the ‘A’ (accumulation increment) inputs change only at low frequency, and the ‘B’ inputs change only when ‘clkn’ is active. In this configuration, the ‘X’ outputs of the first stage are settled before the logic cascoded with ‘C’ and ‘clkp’ in the front-end of the second stage becomes active. As a result, the added stage has no impact on propagation delay. The second stage of the sum has only three cascaded levels as opposed to four in the previous design, so it should actually have a slightly faster overall propagation delay. While the new sum
Figure 3.24: Layout view of the ACCV3 4-bit accumulator test chip. The test chip includes DAC output circuitry and is 1725 µm by 1025 µm. The 4-bit accumulator occupies an area of 510 µm by 575 µm.

circuit is faster than in the previous designs, the single-level parallel-gated carry circuit dominates the critical timing path. The net effect is that both designs have roughly the same maximum clock frequency, but the power is reduced in ACCV3.

The ACCV3 4-bit accumulator has a similar layout topology to the ACCV1 and ACCV2 designs in Section 3.2.3 and Section 3.2.4, except that the accumulator circuit is rotated by 90 degrees. The test circuit also uses a redesigned DAC circuit that has improved linearity. The test circuit layout spreads out the transistors in order to minimize the thermal density of the design. A layout view of ACCV3 test circuit is shown in Figure 3.24. The test chip includes all of the bond pads and the DAC circuitry. The 4-bit accumulator occupies an area of 510 µm by 575 µm. The layout view is shown because a microphotograph of ACCV3 was not captured.
Table 3.9: Yield for the ACCV3 accumulator test circuit in fabrication run TC7. A device is considered to pass if it operates correctly above a 24 GHz clock rate.

3.2.5.1 Simulation Results

The ACCV3 accumulator is simulated to operate up to a 40 GHz clock frequency, and the simulation output is shown in Figure 3.25. The three-level series-gated sum circuit allows for a reduction in the power supply voltage and a significant reduction in power consumption. Some of the design data concerning the ACCV3 design was lost after tapeout, so a full reporting of a simulated power breakdown for the chip is not available. However, measured power data is given in the next section.

3.2.5.2 Measurement Results

The ACCV3 accumulator test circuit was tested on-wafer using using the test setup described in Section 3.2.2. In this fabrication run (TC7), the circuit yield was poor, and there were only four operational ACCV3 accumulators out of 37 tested, as shown in Table 3.9. Similarly, the yield was low for other baseline designs that were included on the wafer in this fabrication run. Approximately one-fourth of the non-yielding circuits were due to electrical shorts in defective capacitors. A device is considered to pass if it operates correctly above a 24 GHz clock rate.

The maximum operating clock frequency for the ACCV3 accumulators was 34 GHz, or 15% slower than expected from simulation. Other baseline designs on the test wafer performed slower by a similar factor. Although the available data is limited by the poor yield, it is likely that this particular wafer and fabrication run were near a “slow” process corner. Output waveforms for ACCV3 were not captured.
Figure 3.25: Simulation of the ACCV3 4-bit accumulator using the single-level parallel-gated carry circuit and the three-level series-gated sum. This simulation includes parasitic extracted capacitors and is shown at the maximum operation frequency of 40 GHz.
The power breakdown for the accumulator components is estimated from a combination of simulation and measured data. The ACCV3 4-bit accumulator consumes a total of 1.97 W of power. The two 2-bit adders each consume 0.59 W of power, the 2-bit register consumes 0.31 W of power, and the clock tree circuitry consumes 0.48 W of power. If the clock tree is divided so that the adders and registers include the power from the clock tree circuitry that they require, the 2-bit adders would use 0.75 W and the 2-bit registers would use 0.47 W.

### 3.2.6 Accumulator ACCV4

The ACCV4 accumulator uses a nearly identical architecture as the ACCV3 accumulator, with the exception of clock tree circuitry. The two designs also use different layout topologies. While the ACCV3 4-bit accumulator uses a layout with spread out transistors to achieve low thermal density, the ACCV4 accumulator uses a more compact layout topology with a reduced transistor to transistor spacing. While this increases the thermal density, it reduces the length of the signal and clock interconnects and the corresponding parasitic capacitances on these interconnects.

The ACCV4 accumulator also has additional clock buffer circuitry as a conservative approach to ensure that the latches have strong clock inputs. A layout view of ACCV4 test circuit is shown in Figure [3.26](#). While it is in the same padframe as the ACCV3 test circuit, the ACCV4 4-bit accumulator circuitry (excluding metal fill and other circuitry) occupies an area of 370 µm by 350 µm. This is a 44% reduction in accumulator area compared to the ACCV3 design. The layout view is shown because a microphotograph of ACCV4 was not captured.
3.2.6.1 Simulation Results

Simulations results indicate that the ACCV4 accumulator has a maximum clock frequency of 43 GHz, as shown in Figure 3.27. The 3 GHz clock frequency improvement over ACCV3 is due to two factors. First, the additional clock buffer circuitry improves the clock signal, which improves latch performance. Second, since the layout is more compact, parasitic capacitance on the signal paths is reduced, which reduces the propagation delay of the signals and increases the maximum clock frequency.

The three-level series-gated sum circuit allows for a reduction in the power supply voltage and a significant reduction in power consumption as was shown by ACCV3, although the higher power clock circuitry in ACCV4 offsets some of the power savings. As with ACCV3, a loss of design data prevents the reporting of a full power breakdown for the chip, but measured power results are included in the next section.
Figure 3.27: Simulation of the ACCV4 4-bit accumulator using the single-level parallel-gated carry circuit, three-level series-gated sum circuit, and additional clock buffer circuitry. This simulation includes parasitic extracted capacitors and is shown at the maximum operation frequency of 43 GHz.
<table>
<thead>
<tr>
<th>Result</th>
<th>Number of Devices</th>
</tr>
</thead>
<tbody>
<tr>
<td>Fail (Short)</td>
<td>6</td>
</tr>
<tr>
<td>Fail (Other)</td>
<td>9</td>
</tr>
<tr>
<td>Pass</td>
<td>1</td>
</tr>
</tbody>
</table>

Table 3.10: Yield for the ACCV4 accumulator test circuit in fabrication run TC7. A device is considered to pass if it operates correctly above a 24 GHz clock rate.

### 3.2.6.2 Measurement Results

The ACCV4 accumulator test circuit was tested using the test setup described in Section [3.2.2](#). This circuit was part of the TC7 fabrication run, which showed poor circuit yield. The ACCV4 accumulator had worse yield than the other designs, with only one operational site out of 16 tested. Defective capacitors accounted for one-third of the non-yielding circuits in this design. The yield for ACCV4 in fabrication run TC7 is shown in Table [3.10](#). A device is considered to pass if it operates correctly above a 24 GHz clock rate.

The single operating ACCV4 accumulator operated at a 35 GHz clock frequency, which is 18.6% slower than expected from simulation. This was comparable to other baseline designs on the test wafer. As stated in Section [3.2.5.2](#), it is likely that this particular wafer and fabrication run were near a “slow” process corner. An output waveform for ACCV4 was not captured.

The power breakdown for the accumulator components is estimated from a combination of simulation and measured data. The 4-bit accumulator consumes a total of 2.54 W of power. The two 2-bit adders each consume 0.59 W of power, the 2-bit register consumes 0.31 W of power, and the clock tree circuitry consumes 1.05 W of power. If the clock tree is divided so that the adders and registers include the power from the clock tree circuitry that they require, the 2-bit adders would use 0.94 W and the 2-bit registers would use 0.66 W.
### Design

<table>
<thead>
<tr>
<th>Design</th>
<th>Maximum Clock Frequency</th>
</tr>
</thead>
<tbody>
<tr>
<td>T. Mathew et al. [27][59]</td>
<td>19 GHz</td>
</tr>
<tr>
<td>Four-Level Series-Gated (ACCV1)</td>
<td>52 GHz</td>
</tr>
<tr>
<td>Single-Level Parallel-Gated (ACCV2)</td>
<td>55 GHz</td>
</tr>
</tbody>
</table>

Table 3.11: Comparison of divide by two carry test circuits. The four-level series-gated and single-level parallel-gated designs in this work exceed the performance of the previous work by more than 2.7 times the clock frequency. Both of the divide by two carry test circuits results are from fabrication run TC6.

#### 3.2.7 Summary of High-Speed 4-bit Accumulators

The four 4-bit accumulator designs establish a comparison to previous work by T. Mathew et al. [27][59], then extend beyond the previous work using modified sum and carry circuit designs. The new accumulators go beyond the previous work by achieving fabricated 4-bit accumulator designs as opposed to accumulator components only. The ACCV3 and ACCV4 designs show that power can be reduced while still achieving high clock frequencies.

Compared to the previous four-level series-gated divide by two carry test circuit [27][59] that operated up to a maximum clock frequency of 19 GHz, the four-level series-gated and single-level parallel-gated divide by two carry test circuits in Section 3.2.3 and Section 3.2.4 are more than 2.7 times faster. A summary of the divide by two carry test circuit results is shown in Table 3.11.

The comparison of the 4-bit accumulators is not straightforward. ACCV2 was measured from fabrication run TC6, while the three other designs have measured results from fabrication run TC7. By comparing other test structures on the TC7 wafer, it was discovered that TC7 was a “slow” fabrication run. As a result, the comparison of the 4-bit accumulator results is not strictly a comparison of design differences only, since there are some process differences embedded in the results. The maximum clock frequency results for ACCV1, ACCV3, and ACCV4 are likely to be lower than the ACCV2 design partially because of the process variations, so the comparison is slightly
<table>
<thead>
<tr>
<th>Design</th>
<th>Fabrication Run</th>
<th>Maximum Clock Frequency</th>
<th>Power Consumption</th>
<th>Speed/Power Ratio</th>
</tr>
</thead>
<tbody>
<tr>
<td>ACCV1</td>
<td>TC7</td>
<td>38 GHz</td>
<td>2.76 W</td>
<td>13.77 GHz/W</td>
</tr>
<tr>
<td>ACCV2</td>
<td>TC6</td>
<td>41 GHz</td>
<td>3.04 W</td>
<td>13.49 GHz/W</td>
</tr>
<tr>
<td>ACCV3</td>
<td>TC7</td>
<td>34 GHz</td>
<td>1.97 W</td>
<td>17.26 GHz/W</td>
</tr>
<tr>
<td>ACCV4</td>
<td>TC7</td>
<td>35 GHz</td>
<td>2.54 W</td>
<td>17.38 GHz/W</td>
</tr>
</tbody>
</table>

Table 3.12: Comparison of 4-bit accumulator circuits. The differences in circuit performance are partially due to design differences and partially due to process variations. Results from test structures indicate that the TC7 fabrication run was slower than the TC6 fabrication run.

flawed. A comparison of the accumulator results is given in Table 3.12. Note that since the previous work by T. Mathew et al. [27, 59] did not construct a working accumulator, it can not be compared to the 4-bit accumulator designs in this thesis.

Even without compensating for the “slow” fabrication run, it is clear that it is possible to reduce power consumption, yet still maintain high-speed operation by moving from a four-level series-gated architecture to an architecture that uses fewer levels. Comparing the results for ACCV2 and ACCV3 from Table 3.12, a 35% reduction in power consumption is achieved with only a 17% reduction in speed. Had the process variation in TC7 been more favorable, the speed difference would have been even smaller. Even comparing ACCV2 to ACCV4, the reduction in power consumption to the reduction in speed is still favorable, with a 16% reduction in power consumption and a 14% reduction in speed. It is interesting to note that while ACCV3 had the best speed/power ratio, the other three designs had nearly identical speed/power ratios. The ACCV4 design shows that the more compact layout works, and does not suffer from the increase in thermal density. This is important, because compact layouts are necessary for extending the accumulator design to larger bit-widths for use in DDS circuits. This more compact layout topology is used in all the subsequent larger designs in this thesis.
3.3 Low Power 8-bit Accumulators

All of the accumulators discussed in Section 3.2 (ACCV1 through ACCV4) were designed with maximum clock frequency as the most important design metric. The high operating frequencies of those accumulators comes at the cost of high power consumption. Even with the modifications to reduce power in ACCV3 and ACCV4, the total power consumption is still quite high. At best, it is nearly 2 W for a 4-bit accumulator. Larger bit-width accumulators that use even more power are needed for DDS designs. This section focuses on low power 8-bit accumulators that are used in DDS circuits.

Low power is achieved by reducing the clock rate of the accumulators. With a target clock rate of around 30 GHz for X-band and Ku-band DDS applications, the extra power needed to achieve accumulators with 40 GHz clock rates is unnecessary for the DDS. The ACCV5 accumulator discussed in Section 3.3.1 has clock rate of 32 GHz. It reduces the power consumption by reducing the maximum operating speed, and it also extends its bit-width to 8 bits for use in a DDS. The 8-bit ACCV6 accumulator reduces power consumption even further by operating at an even lower 13 GHz clock rate. While this DDS is below the target of 30 GHz, it is useful for showing the effect a reduced frequency has on power consumption and other accumulator architectures that are possible. The DDS circuits that these 8-bit accumulators are integrated into will be discussed in further detail in Chapter 4.

3.3.1 Accumulator ACCV5

In order to reduce the power consumption in the accumulator, it is important to consider that much of the accumulator circuitry, particularly latches, requires a power supply that supports two diode drops plus the overhead for voltage swing and the current source. The three-level series-gated sum circuit, however, requires a supply that supports three diode drops plus the overhead for voltage swing and the current source. Thus, it constrains the voltage supply to a voltage higher than necessary for the latches, and it
Figure 3.28: Two-level parallel-gated sum circuit and separate latch circuit.

does not allow for further power reduction. By separating the logic gate and the latch, the three-level series-gated sum circuit becomes a two-level parallel-gated sum circuit, as shown in Figure 3.28. This sum circuit is simply two XOR gates followed by a latch. As was the case in ACCV3 and ACCV4 (described in Section 3.2.5 and Section 3.2.6), the first XOR gate is driven by the accumulation increment and the sum from the previous state, so its Xp/Xn output settles shortly after the clock transition. Therefore, the overall propagation delay of the sum circuit is dominated by the propagation delay of the second XOR gate, which is dependent on the carry input. By using the two-level parallel-gated sum circuit, one of the diode drops can be removed from the power supply requirement, reducing it from -4.6 V to -3.8 V. Since the carry circuit is already parallel-gated, it remains unchanged in the ACCV5 accumulator.

While low power accumulators are desired for DDS designs, accumulators with larger bit-widths are also necessary. Unlike the aforementioned 4-bit accumulators, the ACCV5 accumulator is extended to 8-bits. Using 2-bit accumulator and 2-bit register blocks, the 8-bit accumulator has the topology shown in Figure 3.29. The dotted boxes in Figure 3.29 partition the 8-bit accumulator into 4-bit accumulator and 4-bit register...
blocks. The 8-bit accumulator is essentially an extension of the pipelined 4-bit accumulator shown in Figure 3.9.

### 3.3.1.1 Simulation Results

The overall propagation delay of the sum circuit is dependent on the carry circuit. In simulation, this propagation delay was 6.35 ps. Coincidentally, the carry circuit also had a simulated propagation delay of 6.35 ps. The latch circuit has a simulated propagation delay of 8.37 ps. The propagation delays were measured from sum, carry, and latch circuits integrated in an accumulator circuit, so the loading and extracted parasitics are accounted for. Since the simulated propagation delays for the sum and carry circuits are equal, the critical path of the accumulator is either through one carry gate (6.35 ps) plus one sum gate (6.35 ps) and two latches (2x8.37 ps), or through two carry gates (2x6.35 ps) and two latches (2x8.37 ps). Using the total propagation delay of either critical path, the propagation delay of the accumulator is expected to be 29.44 ps. The critical path propagation delay represents the shortest clock period, or the maximum operating frequency. In this case, the maximum estimated operating frequency of the accumulator is approximately 34 GHz. This is determined from the worst case transitions. It should be noted that skew in the data paths prevents the design from improving the speed by tuning the clock delay. A simulation of the ACCV5 8-bit accumulator at the maximum 34 GHz clock frequency is shown in Figure 3.30. The analysis of the accumulator speed is slightly optimistic. While it includes loading and parasitic capacitances internal to the accumulator, it ignores additional loading and parasitic capacitance from circuitry that the accumulator would drive in practice.

Unlike the previously discussed designs, the sum and carry circuits in ACCV5 are not cascaded with latches. While this reduces the maximum clock frequency, it allows for the elimination of two diode drops from the voltage supply compared to the ACCV1 and ACCV2 designs in Section 3.2.3 and Section 3.2.4 and for the elimination
Figure 3.29: Block diagram of the pipelined 8-bit accumulator using 2-bit adders and 2-bit registers. The dotted boxes partition the accumulator into 4-bit accumulator and 4-bit register blocks.
Figure 3.30: Simulation of the ACCV5 8-bit accumulator. This simulation includes parasitic extracted capacitors and is shown at the maximum operation frequency of 34 GHz. This simulation is optimistic, because it ignores the additional loading and parasitic capacitance from interconnects and circuitry that would be connected to the output of the accumulator in practice.
of one diode drop from the voltage supply compared to the ACCV3 and ACCV4 designs in Section 3.2.5 and Section 3.2.6.

3.3.1.2 Measurement Results

Unlike the previous accumulators, ACCV5 is not implemented inside an accumulator test circuit with a DAC output. Instead, it is implemented inside of a whole DDS circuit. Because of this, no direct output plots of the accumulator are captured. However, the results from the DDS testing described in Section 4.4 showed that the accumulator portion of the circuit operates up to a 32 GHz clock frequency. This is 2 GHz lower than expected from simulation, but this is reasonable considering that the simulated circuit only included the accumulator and did not factor in the DDS components connected to the accumulator output. These components add extra loading and parasitic capacitance to the accumulator feedback interconnects, lowering the operating frequency. The fabricated ACCV5 accumulator was measured up to a 32 GHz maximum clock frequency, so it exceeded the 30 GHz clock frequency requirement for an X-band DDS. Including the clock tree circuitry with the components that it drives, the ACCV5 8-bit accumulator uses 4.89 W, with two 4-bit accumulators at 1.84 W each and a 4-bit register at 1.21 W. Comparing only the 4-bit accumulator portion, the ACCV5 design reduces power by more than 10%. When the accumulators are extended to 8-bits, the difference is even larger. A comparison of the accumulators is presented in Section 3.5. The ACCV5 accumulator is reported as part of a DDS design in [51].

3.3.2 Accumulator ACCV6

The ACCV5 accumulator design in Section 3.3.1 illustrates that power can be reduced by relaxing the clock frequency requirements. If the frequency requirements are relaxed even further to 10 GHz, more architecture changes can be made to facilitate
power reduction. This is the goal of the ACCV6 accumulator. In the previously discussed accumulators, the register and latch circuits consume a large portion of the total accumulator power because of the high degree of pipelining that is necessary to achieve high frequency operation. At lower frequencies, less pipelining is needed and many of the registers and latches can be eliminated. This reduces power not only by eliminating register and latch circuits, but by eliminating the clock tree circuitry that is required by the clock inputs of these circuits.

ACCV6 is an 8-bit accumulator and is shown in Figure 3.31. Instead of using a latch after every logical operation, as in the previous designs, the ACCV6 accumulator uses only one stage of pipelining for the full 8-bit accumulation operation. This architecture resembles a carry ripple adder, with the critical timing path through the string of carry circuits. The full adder shown in Figure 3.32 is the basic building block for this accumulator. It uses a single-level carry circuit with an emitter follower output, but without an output latch as shown in Figure 3.33. Likewise, the sum circuit shown in Figure 3.34 is based on the two-level sum circuit in Figure 3.28. It is integrated with emitter followers so it can drive the pipeline register. Since the ACCV6 architecture eliminates the 2-bit and 4-bit register blocks along with the latches associated with the carry circuits inside of the accumulators, large power savings are achieved.

3.3.2.1 Simulation Results

In the ACCV6 accumulator, the timing is dominated by the propagation delays through the carry circuits. The worst case propagation delay in the accumulator occurs when the carry signal propagates though all of the carry circuits and causes the MSB of the accumulator to change state. The worst case propagation delay that determines the maximum clock frequency is the sum of the propagation delay through seven carry circuits, one sum circuit, and one latch. The propagation delays of the circuits are simulated from the layout of the accumulator with parasitic capacitances included. The
Figure 3.31: Schematic of the ACCV6 8-bit accumulator. It has an architecture similar to a carry ripple adder.

Figure 3.32: Schematic of the ACCV6 full adder block.
Figure 3.33: Schematic of the ACCV6 carry circuit. It is similar to the other single-level parallel-gated carry circuits, except that it is not followed by a latch that recovers full differential. Instead it has an emitter follower so that it can drive subsequent carry and sum circuits.
Figure 3.34: Schematic of the ACCV6 sum circuit. It is similar to other two-level parallel-gated sum circuits, except that it has an emitter follower so that it can drive a register circuit.
propagation delay of the carry circuit is 8.11 ps, for the sum circuit it is 7.58 ps, and for the latch it is 10.41 ps. This represents a total worst case propagation delay of 74.76 ps, for a maximum clock frequency of 13.4 GHz. To determine the maximum clock frequency of the accumulator with parasitic capacitances in simulation, the clock is adjusted in 1 GHz increments. The ACCV6 8-bit accumulator simulated up to a maximum clock frequency of 13 GHz. This simulation output is shown in Figure 3.35.

3.3.2.2 Measurement Results

The ACCV6 accumulator was integrated as part of a DDS circuit described further in Section 4.5, so no direct measurements from the accumulator are available. However, the accumulator is expected to dominate the critical timing of the DDS. Since the full DDS circuit worked up to a 13 GHz clock frequency, the maximum operating frequency of the accumulator is expected to be 13 GHz. This correlates with the results from simulation.

Including all of the clock tree circuitry, the ACCV6 8-bit accumulator consumes 2.13 W of power. Since this accumulator has a different architecture than the previous designs and is not based on 2-bit adder blocks, there is no direct measurement for a 4-bit accumulator. For comparison purposes, half of the power of the 8-bit accumulator, or 1.07 W, can be used as a power figure for the 4-bit accumulator. Since ACCV6 has a different architecture from the other designs, it is the only design for which this method of estimation is valid. In the other designs, extra pipeline registers are needed to extend a 4-bit accumulator design to an 8-bit accumulator design. Also, an extension of the ACCV6 accumulator to larger bit-widths is based on an 8-bit accumulator building block instead of 2-bit or 4-bit accumulator blocks. The ACCV6 accumulator is reported as part of a DDS design in [62].
Figure 3.35: Simulation of the ACCV5 8-bit accumulator. This simulation includes parasitic extracted capacitors and is shown at the maximum operation frequency of 13 GHz.
3.3.3 Summary of Low Power 8-bit Accumulators

While both ACCV5 and ACCV6, which are described in Section 3.3.1 and Section 3.3.2, are 8-bit low power accumulators, they have very different characteristics. ACCV5 uses an architecture that is more similar to the 4-bit accumulators of Section 3.2, based on 2-bit accumulator blocks. Unlike the 4-bit accumulators, however, ACCV5 separates the logic and latch circuitry in order to reduce the voltage supply and the corresponding power consumption. Since ACCV5 is extended to 8-bits, there are more register circuits, so the supply voltage reduction has a more significant impact on power consumption.

ACCV6 uses an architecture that is vastly different from the 4-bit accumulators described in Section 3.2 and ACCV5. Like ACCV5, ACCV6 is an 8-bit accumulator with the logic and latches separated, for a reduced supply voltage compared to the 4-bit accumulator designs. Instead of using the 2-bit accumulator as the basic building block, ACCV6 instead uses the full 8-bit accumulator as a basic building block. This eliminates many of the pipeline registers that are inherent in the other designs. While this reduces the maximum clock frequency, it also greatly reduces the power consumption.

A comparison of the ACCV5 and ACCV6 8-bit accumulators is shown in Table 3.13. The speed-power ratios for both designs are roughly equal, so for the 8-bit accumulator speed and power scale by about the same amount. If these designs are extended to larger bit-widths, the ACCV6 design would have a better speed/power ratio than the ACCV5 design, because it would need fewer registers.

3.4 Reduced Power Accumulator Experiments

As accumulator bit-widths are extended beyond 8-bits, further reductions in power consumption may be necessary. Maintaining high-speed operation is also necessary, so other techniques for power reduction must be explored. These other techniques
### Table 3.13: Comparison of 8-bit accumulator circuits.

<table>
<thead>
<tr>
<th>Design</th>
<th>Maximum Clock Frequency</th>
<th>Power Consumption</th>
<th>Speed/Power Ratio</th>
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<tr>
<td>ACCV5 [51]</td>
<td>32 GHz</td>
<td>4.89 W</td>
<td>6.54 GHz/W</td>
</tr>
<tr>
<td>ACCV6 [62]</td>
<td>13 GHz</td>
<td>2.13 W</td>
<td>6.10 GHz/W</td>
</tr>
</tbody>
</table>

The ACCV6 design uses less pipelining than the ACCV5 design, so it operates at a lower clock frequency, but it also has lower power consumption. Overall, speed and power scale about the same when comparing these two 8-bit designs.

include triple-tail circuitry and resistor-only current sources. Both approaches allow for lower supply voltages and a corresponding lower power consumption.

In Section 3.4.1, the use of the triple-tail approach [56, 57] is examined. The triple-tail approach allows for the reduction of the power supply voltage by an additional diode drop compared to the ACCV5 design described in Section 3.3.1. While a promising approach, the triple-tail circuitry doesn’t have the performance necessary to maintain high-speed operation. Also, despite the reduction in the supply voltage, the triple-tail approach requires more current sources and ends up consuming slightly more power than other approaches. Since the results in simulation were poor, this design was not fabricated.

In Section 3.4.2, the ACCV5 design described in Section 3.3.1 is modified to use resistor-only current sources instead of the current mirrors described in Section 2.6. This allows for a reduction in the power supply voltage by the full current source overhead, which includes a diode drop and the degeneration resistor voltage drop. Using resistor-only current sources significantly reduces the power consumption while maintaining high-speed operation, but it is potentially risky, since the current source output resistance is reduced from 3.7 kΩ, as described in Section 2.7, to around 50 Ω. At the time of the publication of this thesis, this design was in fabrication.
3.4.1 Accumulator with Triple-Tail Circuitry

In an attempt to achieve speeds similar to ACCV5 while also reducing power consumption, an accumulator using the triple-tail approach [56, 57] is examined. The triple-tail approach aims to reduce power consumption by reducing the supply voltage. By eliminating stacked gates, and moving the lower differential pair transistors up to the upper differential pairs as a third tail, the supply voltage can be reduced by a diode drop from -3.8 V to -3.0 V.

A schematic of a triple-tail latch circuit is shown in Figure [3.36]. In the latch, the clock differential pair is used as the third tail for the logic input differential pairs. For proper switching, the logic and clock inputs must be offset by half of the differential voltage. The logic levels for the triple-tail latch are illustrated in Figure [3.37]. The alignment of the voltage levels of the inputs is important, since the triple-tail circuits will operate only if there is a half voltage swing difference between the inputs.
Figure 3.37: Voltage levels in the triple-tail latch. The voltage levels are compatible with other triple-tail gates, where the clock signal is substituted for the logic signal on the third tail.

Using the latch circuit as an example, when ‘clkn’ is high, the third tail ‘clkn’ transistor has a higher input voltage than either of the logic inputs on the left side of the latch and a majority of the current goes through that transistor. On the right side of the circuit, the ‘clkp’ transistor is in the logic low state, and although the voltage levels are offset, it is still lower than the logic high of the A input. In effect, the clock transistor clamps the current on the left side of the circuit, so only the right side of the circuit determines the output.

The triple-tail sum circuit is modified from the ACCV5 sum circuit, and it is shown in Figure 3.38. The move to a triple-tail approach requires two extra current sources and two extra emitter follower circuits. The extra emitter followers are needed to ensure that the triple-tail inputs are spaced properly.

The triple-tail carry circuit is shown in Figure 3.39. Compared to the ACCV5 carry circuit, emitter followers are added so that the voltage going into the latch circuit is at the proper level. Since the carry has a non-standard voltage swing, there may be some performance degradation when it is integrated with the latch circuit.

The triple-tail circuits are compared to their ACCV5 counterparts by both speed and power consumption in Table 3.14. The speed figures are determined by simulations with the circuits configured as clock divide by two circuits. In the divide by two
Figure 3.38: Schematic of the triple-tail sum circuit.
Simulations, the carry and sum circuits also contain latch circuits. The simulations are run with schematics only, without any parasitic capacitances for both ACCV5 and the triple-tail design. Schematic simulations are used so that a layout of the ACCV6 triple-tail components is not needed. The ACCV5 components must also be simulated with schematics only so that the results can be compared. In the power measurements, the triple-tail circuits have a -3.0 V supply voltage, and the ACCV5 circuits have a -3.8 V supply voltage.

From Table 3.14, there is little difference in the triple-tail and ACCV5 latch circuits. The power difference is negligible, and the triple-tail circuit is 4.4% slower. It is possible that the speed difference might be reduced in layout, because the clock signal might drive the triple-tail circuits better since it only goes through a single emitter follower, instead of the double emitter follower that is necessary to obtain the correct voltage level in the ACCV5 circuit. The sum circuits have similar performance, but the
### Triple-tail Speed and Power Comparison

<table>
<thead>
<tr>
<th>Circuit</th>
<th>Triple-tail Speed</th>
<th>Triple-tail Power</th>
<th>ACCV5 Speed</th>
<th>ACCV5 Power</th>
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<td>91 GHz</td>
<td>53.8 mW</td>
</tr>
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<td>Sum</td>
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<td>107 mW</td>
<td>58 GHz</td>
<td>34.0 mW</td>
</tr>
<tr>
<td>Carry</td>
<td>59 GHz</td>
<td>56.7 mW</td>
<td>74 GHz</td>
<td>38.0 mW</td>
</tr>
</tbody>
</table>

Table 3.14: Comparison of speed and power between triple-tail and ACCV5 accumulator components. In the comparison, the components are configured as divide by two circuits and the simulations are of the schematics without extracted parasitics.

The triple-tail circuit uses more than three times as much power. Although the supply voltage is reduced from the ACCV5 circuits to the triple-tail circuits, the extra current sources needed to achieve proper logic voltage levels in the sum circuit increase the total power. The triple-tail carry circuit requires emitter followers that are not needed in the ACCV5 carry, so the power is increased. The triple-tail carry circuit is also much slower than the ACCV5 carry. The carry circuit uses a non-standard voltage level as described in Section 3.2.4, but the triple-tail circuits need to have close alignment between the inputs for proper operation. The non-standard voltage level of the carry circuit does not match ideally with the latch, so performance is reduced.

The triple-tail circuits end up requiring more power and they operate at a lower speed than the ACCV5 circuits. Also, the alignment required by the gate inputs makes the triple-tail circuitry risky to implement. Since the triple-tail circuits did not show any improvement over the ACCV5 circuits, they were not implemented into a full accumulator design or fabricated.

### 3.4.2 Accumulator with Resistor-Only Current Sources

Before the discussion of the accumulator with resistor-only current sources can begin, a baseline for comparison must first be established. In DDS circuits, large bit-width accumulators allow for fine frequency resolution. However, accumulator outputs are often truncated so that the full bit-width of the accumulator is not used for phase...
Figure 3.40: Schematic of a 12-bit accumulator truncated to 9-bits of output. Since the accumulator output is truncated, registers are eliminated to save power.

conversion. When this approach is taken, registers can be eliminated from the accumulator to save power. An example of this is a 12-bit accumulator that is truncated to a 9-bit output. This circuit is shown in Figure 3.40. The ACCV5 accumulator described in Section 3.3.1 is extended to the architecture illustrated in Figure 3.40 for use as a baseline for comparison. In simulation, this baseline accumulator operates up to a maximum clock frequency of 34 GHz, and it consumes 8.58 W of power. Note that this power figure also includes the clock tree circuitry.

The baseline for comparison is modified to use resistor-only current sources. The basic circuits, such as the sum, carry and latch are modified from the circuits used in Section 3.3.1, so that all of the same values of current are used for every circuit. This provides a good comparison of the resistor-only current source design technique, since the only change between the new design and the baseline is current source.

While removing the current source should allow for a voltage supply reduction of a diode drop, plus the degeneration resistor voltage drop, reducing the voltage supply by this amount would result in very small resistor values for current sources. To allow
for more overhead, the supply voltage is reduced to 3.0 V. With this supply voltage, the smallest current source resistors are 50 Ω. This new accumulator simulates up to a maximum clock frequency of 33 GHz and consumes 6.06 W of power.

The use of the resistor-only current sources results in a 29.4% reduction in power, with only a 2.9% reduction in maximum operating frequency. While the new design simulations show a large improvement over the baseline design, the use of resistor-only current sources may be risky because of a reduced current source impedance, and could potentially have problems when fabricated. The original current sources described in Section 2.6 had a much higher impedance of 3.7 kΩ compared to the worst case of 50 Ω for the resistor-only current sources. Variations in the voltage across the current source resistor will have a much larger impact on the current than in the baseline that uses current mirrors. There is a risk that a change in the voltage across the current source resistor could reduce the current of a gate, so that it no longer has a sufficient differential voltage to drive subsequent gates. Despite the potential risks, the resistor-only current source accumulator will be fabricated. The measured results are expected to indicate the viability of the design. At the time of the publication of this thesis, this accumulator design was still in fabrication.

3.4.3 Summary of Reduced Power Accumulator Experiments

Two design approaches, triple-tail and resistor-only current sources, were attempted to reduce power and maintain high-speed operation. While previously reported use of the triple-tail approach [56, 57] showed promise, it did not work out well in accumulator simulations. Not only was the power consumption higher than a comparable circuits from the ACCV5 design, but maximum clock frequency was reduced. Since the triple-tail approach showed no benefits in simulation, it was not fabricated.

The resistor-only current source accumulator had much better simulation results. Compared to a baseline design, it operated only 2.9% slower while consuming 29.4%
less power in simulation. The reduction in the current source impedance from 3.7 kΩ to 50 Ω is concerning and considered risky, since the gate currents may vary greatly as the voltage across the current source resistor changes. At the time of publication, these circuits are not back from fabrication, so the level of risk and the measured circuit performance is unknown.

3.5 Summary of Accumulators

As shown in the preceding sections, there are many possible methods for implementing accumulator circuits. In Section 3.2, 4-bit pipelined accumulators that had logic merged with latches were presented. The ACCV1 accumulator in Section 3.2.3 was based on the work by T. Mathew et al. [27, 59]. It was used as a baseline design for comparison purposes. In Section 3.2.4, an alternative carry circuit was used in the ACCV2 accumulator. This single-level parallel-gated carry circuit took a step towards reduced power consumption. On its own, the carry circuit did not lead to reduced power consumption, because the voltage supply was constrained by the four-level series-gated sum circuit. It did, however, provide a proof of concept for the new carry circuit design. In Section 3.2.5 a three-level series-gated sum circuit was introduced for use in the ACCV3 accumulator. Combined with the single-level parallel-gated carry circuit, a reduction in the voltage supply and the corresponding power achieved in ACCV3. The ACCV4 accumulator in Section 3.2.6 was modified from ACCV3 to have both a smaller layout and increased drive in the clock tree circuitry. Using a smaller layout increases the thermal density, but it decreases the parasitics on interconnects and shortens clock distribution interconnections. As a conservative approach, the power in the clock tree was also increased to ensure that the pipeline latches receive strong signals.

Compared to the previous work by T. Mathew et al. [27, 59], ACCV1 and ACCV2 showed increased performance for the divide by two carry test circuit. Performance was increased from a maximum clock frequency of 19 GHz in the previous
work to 52 GHz in ACCV1 and 55 GHz in ACCV2. The 4-bit accumulators in this work also improved upon the previous work, by yielding operating accumulators. In the previous work, only components of accumulators were operational.

In Section 3.3.1, the ACCV5 accumulator uses similar circuitry and layout to the ACCV4 design. Instead of merging the logic and the latches, these elements are separated. While this leads to a slightly lower operating frequency, it also results in lower power consumption because the supply voltage is reduced. The ACCV5 accumulator is also extended to 8-bits and integrated in a DDS test circuit. The DDS circuit (DDSV1) is described in further detail in Section 4.4 and is reported in [51].

The ACCV6 accumulator in Section 3.3.2 is an 8-bit design that uses a different approach from the previous designs. It is designed for a lower operating frequency, so it does not use any intermediate pipeline stages. Eliminating the pipelining greatly reduces the power consumption. The ACCV6 accumulator is also integrated in a DDS test circuit. The DDS circuit (DDSV2) that uses ACCV6 as a phase accumulator is described in further detail in Section 4.5 and is reported in [62].

Work was also undertaken to further reduce power consumption while maintaining high clock frequencies in Section 3.4. This is necessary for further evolutions of DDS designs. The triple-tail circuit in Section 3.4.1 ended up with worse performance and more power consumption than baseline designs, so the approach was abandoned. The resistor-only current source approach in Section 3.4.2 showed promise since it had significantly reduced power consumption with minimal impact on high speed performance. The design may be risky, however, since the current source impedance is greatly reduced. The resistor-only current source approach is in fabrication, so measurement results are not available.

Since the accumulator sizes are different, a comparison of the 4-bit accumulators from Section 3.2 and the 8-bit accumulators in Section 3.3 is not direct. In comparing the 8-bit accumulators to the 4-bit accumulators, the ACCV5 8-bit accumulator has a
<table>
<thead>
<tr>
<th>Design</th>
<th>Power (W)</th>
<th>Speed (GHz)</th>
<th>4-bit Speed/Power Ratio (GHz/W)</th>
</tr>
</thead>
<tbody>
<tr>
<td>ACCV1</td>
<td>2.76</td>
<td>38</td>
<td>13.76</td>
</tr>
<tr>
<td>ACCV2</td>
<td>3.04</td>
<td>41</td>
<td>13.49</td>
</tr>
<tr>
<td>ACCV3</td>
<td>1.97</td>
<td>34</td>
<td>17.25</td>
</tr>
<tr>
<td>ACCV4</td>
<td>2.54</td>
<td>35</td>
<td>13.78</td>
</tr>
<tr>
<td>ACCV5</td>
<td>1.84</td>
<td>32</td>
<td>17.39</td>
</tr>
<tr>
<td>ACCV6</td>
<td>1.07a</td>
<td>13</td>
<td>12.15a</td>
</tr>
</tbody>
</table>

a No 4-bit accumulator block exists for ACCV6, so the power is estimated from the 8-bit accumulator results.

Table 3.15: Speed and power comparison for 4-bit accumulator designs.

The designs are also compared in terms of 8-bit accumulators in Table[3.16]. For the 4-bit accumulator designs (ACCV1 through ACCV4), the 8-bit accumulator power consumption is estimated by the circuitry needed to build an 8-bit accumulator. For these designs, the 8-bit accumulator power is estimated as the power from two 4-bit accumulators and four 2-bit registers. Note that the combination of four 2-bit registers is needed to construct a 4-bit register. The power figures used in these estimates include the clock tree circuitry required by each block, so the extension to 8-bit accumulators includes this additional clock tree overhead.

The ACCV5 design has the best speed/power ratio in both the 4-bit accumulator comparison in Table[3.15] and 8-bit accumulator comparison in Table[3.16] The
reduction in maximum operating frequency incurred by separating the merged logic and latches is compensated for by the proportionally larger reduction in power consumption due to the lower supply voltage. This benefit is more apparent in the 8-bit accumulator speed/power ratios, since the pipeline buffering registers become a dominant portion of the total power consumption as the bit-width increases. Since the ACCV1 through ACCV4 designs have a higher supply voltage than the ACCV5 design, their speed/power ratios suffer in larger bit-width designs because of the extra power that is essentially wasted in the registers.

The impact of increased register power is further illustrated by the ACCV6 design. It has the worst speed/power ratio in the 4-bit accumulator comparison in Table 3.15, yet it has the second best speed/power ratio in the 8-bit accumulator comparison in Table 3.16. Unlike the other accumulator designs, the ACCV6 design eliminates the intermediate pipeline stages, so it requires fewer pipeline registers as the bit-width increases. The reduction in the number of registers leads to a speed/power ratio that becomes comparably better from the 4-bit comparison to the 8-bit comparison. If the design comparison were extended to larger bit-widths, the ACCV6 design speed/power ratio would overtake all other designs including ACCV5, because it would use much less power from registers. Since pipelining performs no function other than to buffer data to increase the operating frequency, it would follow that the optimal design for a specific frequency would contain the least amount of pipelining necessary. This is accomplished by reducing the number of registers to the minimum necessary to operate at a specific frequency. As the bit-width of the design increases, the lower relative power in designs with the minimum number of registers would become more apparent.

To illustrate this point, Table 3.17 shows a hypothetical extension of the accumulators to a 16-bit bit-width. Unlike the accumulator in Section 3.4.2, the 16-bit extension does not have any truncation. For these designs, the impact of differences in the register
Table 3.17: Speed and power comparison for accumulators extended to 16-bit bit-widths.

<table>
<thead>
<tr>
<th>Design</th>
<th>Power (W)</th>
<th>Speed (GHz)</th>
<th>8-bit Speed/Power Ratio (GHz/W)</th>
</tr>
</thead>
<tbody>
<tr>
<td>ACCV1</td>
<td>29.28</td>
<td>38</td>
<td>1.30</td>
</tr>
<tr>
<td>ACCV2</td>
<td>30.40</td>
<td>41</td>
<td>1.35</td>
</tr>
<tr>
<td>ACCV3</td>
<td>19.16</td>
<td>34</td>
<td>1.77</td>
</tr>
<tr>
<td>ACCV4</td>
<td>26.00</td>
<td>35</td>
<td>1.35</td>
</tr>
<tr>
<td>ACCV5</td>
<td>14.62</td>
<td>32</td>
<td>2.19</td>
</tr>
<tr>
<td>ACCV6</td>
<td>5.14</td>
<td>13</td>
<td>2.53</td>
</tr>
</tbody>
</table>

Power becomes very apparent. The power consumption of the 4-bit designs from Section 3.2 greatly increase to levels that are not feasible for fabrication. The speed/power ratios of the 4-bit designs are also worse than the 8-bit designs from Section 3.3. ACCV6 has the best speed/power ratio in this comparison, because it use a lot less register circuits than any of the other designs. It also has a power consumption of 5.14 W, which is reasonable for a fabricated design.

The results of the accumulators lead to some conclusions on choosing the proper design point for a high-speed accumulator design. In general, pipelining is double-edged sword. While it is necessary to achieve high-speed performance, it uses a large amount of power as bit-widths are increased. This is apparent from the extension of the designs to 16-bits in Table 3.17. Accumulators should be designed so that the minimum amount of pipelining necessary to achieve a desired operating frequency is used. This will help to maximize the speed/power ratio. It should be noted that this is valid when the bit-width of an accumulator is relatively large. For small bit-widths, this may be obscured, and there may be cases where more speed/power ratio is higher for designs with more pipelining if the registers are not a significant portion of the circuit. It is also important to keep total power consumption in mind when extending the bit-widths of designs. A 16-bit accumulator that consumes 30 W of power is not feasible for portable devices, because power consumption is simply too high. The 14.62 W and 5.14 W versions that
are extended from ACCV5 and ACCV6 are much more reasonable alternatives that still show high-speed performance.

Ultimately, the bit-widths of accumulators need to be extended for use in DDS circuits. Not only must the power consumption of the accumulator be considered, but there will also be additional power consumption from the phase conversion and DAC circuitry needed to make a DDS. In the next chapter, two DDS circuits will be discussed. The first design will integrate the ACCV5 accumulator described in Section 3.3.1 into a DDS and the second design will integrate the ACCV6 accumulator described in Section 3.3.2 into a DDS.
CHAPTER 4

Direct Digital Synthesizers

As InP HBT technologies improve and allow for faster digital and mixed-signal circuits, more applications become viable. One application of growing interest is frequency synthesis in the microwave range. Synthesized microwave frequencies are useful in areas such as communications, test equipment, and radar. Compared to traditional methods of microwave frequency generation, the use of high-speed digital logic and InP technology allows for improvements in terms of bandwidth, power, and size of the system.

Several methods of frequency synthesis in the microwave frequency range have been recently reported. Examples include optical methods \[63\], fractional-N methods \[64\], resonators \[65\], and direct digital synthesis. While optical methods have been shown to generate frequencies over 60 GHz \[63\], they require multiple lasers for operation. Fractional-N methods have similar performance to direct digital synthesizers, but they sacrifice agility, or frequency switching time, to improve spurious free dynamic range \[66\]. Resonators are able to generate very pure output signals, but only over very small or fixed frequency ranges. Direct digital synthesizers are able to deliver wide band, high frequency, agile signals with high SFDR.

4.1 DDS Architecture

The basic DDS architecture was first reported by J. Tierney et al. \[67\] over thirty years ago, and it has a fairly simple structure. The DDS block diagram shown in Figure 4.1 consists of a phase accumulator, a phase converter, a digital to analog converter, and a low pass filter (LPF). Conceptually, the DDS operation is simple. The phase accumulator tracks the phase of a sine wave on the unit circle by using a clocked binary adder. The phase converter translates phase information from the accumulator into a
digital representation of the desired output waveform. The DAC then converts the digital output of the phase converter into an analog output. The LPF eliminates unwanted high frequency spectral components from the analog output.

The accumulator has a modular output over the range \([0, 2^N-1]\). This modular output is convenient for mapping the accumulator output to a unit circle representing the phase of a sine wave. This mapping can be represented by

\[
\theta = \frac{(2n + 1)\pi}{2^{N+1}},
\]  

(4.1)

where \(n\) is output of the accumulator. This particular mapping does not include the phases 0, \(\pi/2\), \(\pi\), and \(3\pi/2\). Instead the discrete phase steps are offset by half a step size, so that \(\theta = \pi/2^{N+1}\) for \(n = 0\). This approach simplifies partitioning the phase into quadrants, which is advantageous for the phase converter. An example of this mapping on the unit circle is shown in Figure 4.2.

The phase converter is typically used to generate a sine or cosine output, although any periodic waveform can be implemented [68]. In this thesis, only sine waves are generated by the DDS circuits. There are several methods for implementing the phase converter including a ROM look up table [67], a CORDIC algorithm [69], and
polynomial approximations [70]. Of these methods, one of the most common is the ROM look up table (LUT). LUTs can be simplified by exploiting the symmetry of the unit circle, particularly the quadrant symmetry described above, and by using a combination of coarse and fine value LUTs to form the digital sine value at a particular phase [67]. While ROM look up tables are a simple method of phase to sine conversion, the implementation of a fast ROM is not always possible or feasible. When this is the case, other conversion methods are necessary. One such method uses the iterative CORDIC technique [71] for conversion in a DDS [69]. The CORDIC technique takes advantage of trigonometric identities and approximations to calculate the value of \( \sin(\theta) \) by using shift and add operations. This method overcomes the potential speed limitations of the ROM by using pipelining in the CORDIC conversion. An alternative to the CORDIC method uses the polynomial expansion approximation [70] of the sine function. Circuitry can be simplified by lowering the order of the polynomial used.
for approximation. Likewise, the order of the approximation polynomial can be adjusted to achieve a desired SFDR. Compared to the CORDIC method, the polynomial approximation method uses fewer pipeline stages and less power.

The digital representation of the sine wave is converted into an analog waveform by the DAC. The DAC is not necessarily on the same chip as the accumulator and phase converter, although it generally is in high performance DDS circuits. The DAC is usually noted as the limiting factor \[72\] for spectral purity performance.

The LPF is often implemented by discrete components off-chip, although it is possible to implement it on-chip if on-chip inductors are available. The LPF is needed to eliminate images above the Nyquist frequency or half of the clock frequency \(f_{clk}/2\). Since the filter has roll-off characteristics and must block frequencies above \(f_{clk}/2\), frequencies slightly below \(f_{clk}/2\) are attenuated. This sets a practical limit for a maximum DDS output of about 40\% of \(f_{clk}\) \[73\]. In many cases, the LPF is not included in the DDS circuit design, but it is a component of the whole DDS system.

### 4.2 Performance Metrics and Design Tradeoffs

The performance of the DDS can be measured by various criteria: frequency resolution, maximum frequency, bandwidth, SFDR, agility, and power consumption.

The frequency resolution \(f_0\) is the smallest incremental difference in DDS output frequencies. This is determined by the clock frequency and the bit-width of the accumulator \(N\) as given by

\[
  f_0 = \frac{f_{clk}}{2^N}. \tag{4.2}
\]

A larger bit-width \(N\) will improve the resolution of the DDS and allow for finer frequency spacing. Higher frequency resolution leads to improved SFDR, at the expense of higher power consumption.
The output frequency \( f_{out} \) of the DDS is controlled by the phase increment of the accumulator \( k \) in multiples of the frequency resolution. It is given by

\[
f_{out} = kf_0 = \frac{k f_{clk}}{2^N}.
\]  

(4.3)

Due to the Nyquist sampling theorem, the maximum value for \( k \) is effectively limited to \( 2^N - 1 \), leading to the maximum output frequency \( f_{out\text{max}} \) given by

\[
f_{out\text{max}} = \frac{2^N - 1}{2^N} f_{clk} = \frac{f_{clk}}{2}.
\]  

(4.4)

As noted in the previous section, the maximum DDS output is typically limited to about 40% of \( f_{clk} \).

The bandwidth is the difference between the minimum and maximum output frequencies. The phase increment \( k \) can have a value of 0, so the minimum output of the DDS is 0 Hz. Thus, the bandwidth is given by \( f_{out\text{max}} \). Since the maximum output frequency and bandwidth are determined by \( f_{clk} \), they can be improved by using fast technologies, such as InP HBTs, and design techniques that derive the maximum benefit from the technology. In a particular technology, increasing bandwidth and maximum frequency generally comes at the expense of increased power consumption.

Spurious-free dynamic range is a measure of the spectral purity of the sine wave produced by the DDS. Ideally, the spectral output should only have energy at the desired frequency. Due to factors described below, there is also energy at undesired frequencies. These are referred to as spurious frequencies or spurs. The spurs are periodic, discrete spectral lines \([66]\). For SFDR measurements, only in-band frequencies \((\leq f_{clk}/2)\) are taken into account since the LPF is expected to filter out frequencies above Nyquist. The SFDR is the ratio of the amplitude of the desired signal \( A_p \) to the amplitude of the
Figure 4.3: Frequency spectrum of a DDS output with the desired frequency $A_p$ and largest spurious frequency $A_s$ labelled.

largest in-band spur $A_s$ [69], given by

$$SFDR = 20 \log \left( \frac{A_p}{A_s} \right). \quad (4.5)$$

An example of the output spectrum of a DDS with $A_p$ and $A_s$ labelled is illustrated in Figure 4.3.

Spurs are generated in each stage of the DDS architecture shown in Figure 4.1. Typically, the output of the accumulator is truncated, so not all $N$ bits feed into the phase converter. The quantization of phase information also introduces spurs [66]. The phase converter introduces spurs from compression and quantization [74]. In most implementations of the phase converter, phase compression methods are used to save ROM space or to simplify/reduce calculations. Quantization arises in all types of phase converter, since the output has finite precision. Spurs in the DAC are caused by glitches from bit switching, the finite slew rate of the DAC, and from quantization [72]. It has been noted
that for a low SFDR, the DAC should have low switching transients and a fast settling time [73]. The LPF can also be a cause of spurs [74].

Agility is a measure of how quickly the DDS can change frequencies. Agility is important in applications where the output frequency must hop frequencies quickly, such as in secure communications systems. The agility is defined as the time between changing the frequency control word \( k \) in Figure 4.1 and when the output is at the desired frequency. Using fewer levels of pipelining will tend to improve agility, but it will also tend to decrease the maximum clock frequency.

Power consumption is determined by a combination of the technology used, the clock frequency, the internal bit-width \( N \), and the amount of pipelining internal to the DDS. Since it is dependent on so many factors, it plays a part in most of the performance tradeoffs. Improving any of the other metrics usually results in a design with increased power consumption. In order to decrease the power consumption, the DDS must usually be run at a lower clock frequency, with less frequency resolution, or with a worse SFDR.

### 4.3 Recent DDS Results

With a large number of conflicting tradeoffs, comparisons among reported DDS circuits can be difficult. In CMOS technologies, many of the reported DDS designs are mainly concerned with the demonstration of various phase conversion techniques [69, 70]. These circuits typically have a high SFDR, but a low clock frequency. Designs that are more comparable to the work in this thesis are done in InP technologies. These designs typically try to achieve high maximum clock rates, and SFDR figures in the 20 dBc to 30 dBc range.

In 2001, an InP DDS was reported by A. Gutierrez-Aitken et al. at TRW [8, 75] that had a maximum clock rate of 9.2 GHz and consumed 15 W of power. The authors did not report the SFDR for the full range of FCWs, but of the two FCWs reported, the
worst SFDR was 30 dBC. It is unknown if there were any other FCWs that had a worse SFDR.

In 2005, K. Elliott at HRL reported an InP DDS with a maximum clock rate of 12 GHz that consumed 8 W of power. Up to 40% of the full range of FCWs, the worst case SFDR was 30 dBC. This particular work was part of HRL’s contribution to the TFAST project, so it is the closest reported comparison to the DDS circuits in this thesis.

In the next two sections, two DDS designs employing different tradeoffs will be discussed. Compared to the previous work, these DDS circuits achieve higher maximum clock frequencies and lower power, but with decreased SFDR. In Section 4.4, the DDSV1 design is intended for high-speed operation. To achieve this goal, the SFDR is sacrificed, and the power consumption is high. Overall performance is up to a 32 GHz clock frequency, with a worst case SFDR over the whole range of FCWs of 21.56 dBC and 9.45 W of power consumption. In Section 4.5, the DDSV2 design is intended for low power operation. This tradeoff impacts high-speed performance. The DDSV2 design has a maximum clock frequency of 13 GHz, a worst case SFDR of 26.67 dBC, and 5.42 W of power consumption.

4.4 Direct Digital Synthesizer DDSV1

Unlike the traditional DDS architecture which consists of a phase accumulator, phase converter, and digital to analog converter, DDSV1 combines the phase converter and DAC into a sine-weighted DAC. The resulting DDS architecture is shown in Figure 4.1. The sine-weighted DAC is also an alternative to the traditional approach of using ROM look-up tables. This approach is similar to the cosine-weighted DAC implemented by A. Gutierrez-Aitken et al. [8], but it also adds a Gilbert multiplier for analog inversion. Since the phase converter circuitry is eliminated, this approach allows for a reduction in circuit complexity and power consumption. The accumulator is 8-bits
Figure 4.4: Block diagram of DDSV1 circuit with the outputs of each stage illustrated.

wide, so the DDSV1 frequency resolution is 1/256 of the clock frequency with 128 steps of frequency control. The DDSV1 design uses the ACCV5 8-bit accumulator described in Section 3.3.1, which was simulated to operate up to a 34 GHz clock frequency. The output of the accumulator is truncated, so that only the five MSBs ($a_4$, $a_3$, $a_2$, $a_1$, $a_0$) are used to generate a full-wave sine output.

### 4.4.1 Sine-Weighted Digital to Analog Converter

The sine-weighted DAC uses the five MSBs from the accumulator to generate a full-wave sine output. The DAC is comprised of a thermometer-coder, a sine-weighted summing junction, and a Gilbert multiplier [76]. In the thermometer-coder, the $a_3$ output from the accumulator complements the three LSBs ($a_2$, $a_1$, $a_0$) to expand the thermometer-coder from a quarter-wave to a half-wave output. As shown in Figure 4.5, the complemented bits are registered to meet timing requirements, then buffered to drive the thermometer-coder logic. The thermometer-coder outputs drive the sine-weighted taps of the summing junction, which has a tap weighting scheme of [3 5 5 4 4 3 2 1].
Figure 4.5: Block diagram of the thermometer-coder portion of the sine-weighted DAC.

The DAC is driven by a thermometer coder, thus the weights sum successively to generate a quarter-wave sine output. The first tap weight of 3 is always enabled to ensure that the summing junction has non-zero outputs for all states. This ensures that distinct positive and negative outputs exist for each state after the analog sum is inverted by the Gilbert multiplier. Since the tap weight of 3 is always enabled, the center step level is equal to 6. The Gilbert multiplier uses the accumulator MSB ($a_4$) as a control signal for inversion, resulting in a full-wave sine output. The sine-weighted summing junction and the Gilbert multiplier are shown in Figure 4.6.

The Gilbert multiplier is shown in Figure 4.7. In this DDS it is used to flip the sign of the output of the DDS. It is controlled by the sign bit ($a_4$). Essentially, it functions as a controllable analog inverter, with $a_4$ as the control and DAC output as the analog input. The Gilbert multiplier is advantageous because it is linear over a relatively wide range [76]. Linearity is particularly important for the output of the DDS, because non-linearities negatively impact the SFDR. The resistor connecting the emitters of the input
signal transistors acts as a degeneration resistor and helps to improve linearity further by introducing negative feedback [55].

4.4.2 Phase Truncation Spurs

Since DDSV1 does not use all of the output bits from the accumulator for phase conversion, spurs due to phase truncation arise. Formulas for determining the location and magnitude of these spurs have been developed by V. Kroupa et al. [77]. The magnitude and locations of the spurs are found to be a function of the accumulator bit-width ($N$), the number of bits used for phase conversion ($W$), the number of bits truncated from the accumulator ($B$), the FCW, and the clock frequency. The SFDR (in dBc) that is due to truncation spurs has an upper bound estimate of

$$SFDR \approx 6W(dBc)$$  \hspace{1cm} (4.6)$$

from the analysis in [77]. The actual SFDR is a function of both $W$ and the normalized output frequency ($f_{output}/f_{clk}$) [77], so the upper bound estimate discards the impact from the normalized output frequency. In reality, the SFDR is reduced by a few dBc for
Figure 4.7: Gilbert multiplier schematic. Used to switch the sign of the DAC output in DDSV1 to achieve a full-wave sine output.

normalized output frequencies approaching 1/2. In the DDSV1 design, $W$ is equal to five, so the upper bound on SFDR is about 30 dBc.

Combining terms from [77], the location of the worst case spur from phase truncation is given by

$$f_{spur} = \frac{f_{clk}}{2^N} \left( FCW \pm 2^W P \right), \quad (4.7)$$

where $P$ is the value of the truncated $B$ LSBs of the frequency control word. $P$ can range from 0 to $2^{B-1}$. Note that when these spurs are outside of the Nyquist band (0 to $f_{clk}/2$), they will be aliased back into band. While Equation [4.7] gives the location of the worst case phase truncation spur, other spurs will be located at harmonics of the desired output and harmonics of the worst case phase truncation spur, as well as intermodulations of these signals. Additional spurs occur due to offsets in the Gilbert multiplier and from non-linearities in the DAC.

DDSV1 truncates the output of the 8-bit accumulator to 5-bits, so it has $N=8$ and $W=5$. For the case when $FCW=1$, it is expected that the worse case phase truncation
spurs (taking aliasing into account) will be located at $f_{clk}/8 \pm f_{clk}/256$. Intuitively, when $FCW=1$, the DAC output only changes every eight clock cycles, so spurs located at a frequency of $f_{clk}/8$ mixed with the desired output of $f_{clk}/256$ are expected.

### 4.4.3 Simulation Results

The DDSV1 test circuit simulated up to a maximum clock frequency of 34 GHz. A time-domain simulation output at 34 GHz with an FCW of 1 is shown in Figure [4.8](#).

At this setting, the output frequency of the DDS is 132.8125 MHz. This output has an SFDR of 29.91 dBc, as shown in Figure [4.9](#).

A time-domain simulation output at 34 GHz with an FCW of 127 is shown in Figure [4.10](#). At this setting, the output frequency of the DDS is 16.8671875 GHz. This output has an SFDR of 27.71 dBc, as shown in Figure [4.11](#).

The DDSV1 design uses separate power supplies for the accumulator and the thermometer coder/DAC sections of the circuit. The 8-bit accumulator portion of the
Figure 4.9: Frequency-domain simulation output of DDSV1 with a 34 GHz clock and FCW=1. This output has a 29.91 dBc SFDR.

Figure 4.10: Time-domain simulation output of DDSV1 with a 34 GHz clock and FCW=1. The output frequency is 16.8671875 GHz.
Figure 4.11: Frequency-domain simulation output of DDSV1 with a 34 GHz clock and FCW=1. This output has a 29.91 dBC SFDR.

circuit uses a -3.8 V power supply, and the thermometer coder/DAC portion uses a -4.5 V power supply. The use of separate supplies allows for lower power consumption, since the accumulator portion of the circuit is operated with a lower supply voltage. A complete power breakdown of DDSV1 is shown in Table 4.1

4.4.4 Measurement Results

The DDSV1 test chip was fabricated as part of the TC8 fabrication run. A microphotograph of the test chip is shown in Figure 4.12. It has 1891 transistors in an area of 2700 µm by 1450 µm. As was the case with the accumulator circuits in Chapter 3, DDSV1 was tested on-wafer. The DDSV1 test setup is shown in Figure 4.13. The DDS clock input is differential, but it is driven single-ended, with the non-driven side connected to ground through a 50 Ω termination. The DDS output is also differential, with one single-ended output driving a spectrum analyzer and the other single-ended output
Table 4.1: Simulated power breakdown for DDSV1.

<table>
<thead>
<tr>
<th>Component</th>
<th>Sub-Component</th>
<th>Current</th>
<th>Voltage</th>
<th>Power</th>
</tr>
</thead>
<tbody>
<tr>
<td>8-bit Accumulator</td>
<td></td>
<td>1419 mA</td>
<td>-3.8 V</td>
<td>5.3922 W</td>
</tr>
<tr>
<td></td>
<td>4-bit Accumulator</td>
<td>533 mA</td>
<td>-3.8 V</td>
<td>2.0254 W</td>
</tr>
<tr>
<td></td>
<td>4-bit Register</td>
<td>353 mA</td>
<td>-3.8 V</td>
<td>1.3414 W</td>
</tr>
<tr>
<td>Input Buffers</td>
<td></td>
<td>125 mA</td>
<td>-3.8 V</td>
<td>475 mW</td>
</tr>
<tr>
<td>DAC</td>
<td></td>
<td>1010 mA</td>
<td>-4.5 V</td>
<td>4.545 W</td>
</tr>
<tr>
<td></td>
<td>Thermocoder</td>
<td>370 mA</td>
<td>-4.5 V</td>
<td>1.665 W</td>
</tr>
<tr>
<td></td>
<td>8-bit Register</td>
<td>319 mA</td>
<td>-4.5 V</td>
<td>1.4355 W</td>
</tr>
<tr>
<td></td>
<td>Gilbert Cell/Output Buff.</td>
<td>77 mA</td>
<td>-4.5 V</td>
<td>346.5 mW</td>
</tr>
<tr>
<td></td>
<td>DAC Driver/Core</td>
<td>147 mA</td>
<td>-4.5 V</td>
<td>661.5 mW</td>
</tr>
<tr>
<td></td>
<td>Other</td>
<td>97 mA</td>
<td>-4.5 V</td>
<td>436.5 mW</td>
</tr>
<tr>
<td>Total</td>
<td></td>
<td></td>
<td></td>
<td>10.4122 W</td>
</tr>
</tbody>
</table>

driving a high-frequency sampling oscilloscope. Having both outputs available allows simultaneous testing of both the time and frequency domain outputs of the DDS.

The DDS operates up to a maximum clock frequency of 32 GHz for all frequency control words. A full sweep of all FCWs at 32 GHz is shown in Figure 4.14. For all SFDR measurements, the SFDR is measured within the full Nyquist bandwidth. The worst case SFDR over the range of FCWs is 21.56 dBc at an FCW of 95, which corresponds to an output frequency of 11.875 GHz. The average SFDR over the whole range of FCWs is 26.95 dBc. The maximum operation frequency is better than expected from simulation, since the DDS operates up to a maximum clock frequency of at 32 GHz instead 28 GHz.

A 125 MHz sine-wave output synthesized from a 32 GHz clock frequency with FCW=1 is shown in Figure 4.15. This output represents the fundamental output frequency \( f_0 \) and frequency resolution of the DDS with a 32 GHz clock. With an output frequency of 125 MHz, the DDS has an SFDR of 31.00 dBc with the largest in-band spur located at 4.125 GHz. The frequency spectrum of this output is shown in Figure 4.16. The worst spurs in the output spectrum are due to the truncation of the phase word from the accumulator to the sine-weighted DAC. They are located at \( f_{clk}/8 \pm f_{clk}/256 \), which
Figure 4.12: Microphotograph of the DDSV1 test chip, with dimensions of 2700 $\mu$m by 1450 $\mu$m. The DDSV1 test chip contains 1891 transistors.

Figure 4.13: DDSV1 test setup.
matches the analysis from Equation 4.7. The largest of the additional spurs occur at harmonics of the worst case spurs. The SFDR magnitude is 31.00 dBc, which is close to the 30 dBc result predicted from Equation 4.6 for a DDS using five bits for phase conversion.

To illustrate high output frequency operation of the DDS, the frequency spectrum of the DDS with FCW=127 is shown in Figure 4.17. At this FCW, the output frequency is 15.875 GHz and the SFDR is 30.44 dBc. The largest in-band spur is located at 12.125 GHz, with a similar magnitude spur at 11.875 GHz. These are truncation spurs located aliased frequencies of what would be expected from Equation 4.7.

Of the 32 DDSV1 sites tested, 8 showed either no output or the incorrect output at a 24 GHz clock frequency. The remaining 24 test sites operated at least a 24 GHz clock frequency, and were considered functional for yield purposes. Thus, the functional yield for DDSV1 in fabrication run TC8 was 75%. Overall, the DDS test chip was
Figure 4.15: Sampling oscilloscope output of DDSV1 with $f_{ck} = 32$ GHz and $f_{out} = 125$ MHz.

Figure 4.16: Frequency spectrum of the DDSV1 output with $f_{ck} = 32$ GHz and $f_{out} = 125$ MHz. The largest spur is located at 4.125 GHz and the SFDR is approximately 31 dBc.
Figure 4.17: Frequency spectrum of the DDSV1 output with $f_{\text{clk}} = 32$ GHz and $f_{\text{out}} = 15.875$ GHz at FCW=127. The largest spur is located at 12.125 GHz and the SFDR is 30.44 dBc.

measured to consume a total power of 9.45 W, which is 9.2% lower than the expected from simulation. Extrapolating from Table 4.1 the 8-bit accumulator uses 4.89 W and the DAC uses 4.12 W. DDSV1 is reported in [51].

### 4.5 Direct Digital Synthesizer DDSV2

The DDSV2 design uses a traditional DDS architecture [67] with a phase accumulator, phase converter, and DAC, as shown in Fig. 4.18. This design uses the ACCV6 8-bit accumulator described in Section 3.3.2. Since the accumulator is 8-bits wide, the fundamental output frequency of the DDS is 1/256 of the input clock frequency, and the DDS has 128 steps of frequency control. The 8-bit accumulator output is truncated to 6-bits for use in the phase conversion logic. From the simulations carried out in Section 3.3.2, it was estimated that the ACCV6 accumulator operates up to a 13 GHz clock frequency.
While using a ROM is a more common approach to phase conversion, a ROM design in this technology had not been completed when this DDS was implemented. ROM-less techniques have been reported in other designs \[75, 47\] and are an option for phase conversion. ROM-less approaches are particularly useful in designs with a narrow bit-width, where the amount of conversion logic necessary is small. In this DDS, logic gates are used for a ROM-less phase conversion. The combinational logic performs a look up table operation and drives the DAC. The coding scheme is discussed in further detail in Section \[4.5.1\]

The DAC has seven inputs and is partitioned into a 3-bit coarse DAC and a 4-bit fine DAC. The fine DAC portion is binary, with bit weights of 8, 4, 2, and 1 units of current. In the 3-bit coarse DAC, each bit is 16 units of current. The 3-bit coarse DAC section is driven by thermometer coded outputs from the phase converter. This is in contrast to typical DACs that internally generate thermometer coded signals. Since the DDS design is monolithic, an approach with thermometer coding outside of the DAC can be used. All of the currents from the fine and coarse DAC are summed together through a resistor. This DAC has 64 discrete outputs and is shown in Figure \[4.19\].
4.5.1 Coding Scheme

The coding scheme used in the LUT attempts to approximate the sine wave as accurately as possible, while reducing the complexity of the conversion logic. Using MATLAB, various coding approaches were tested, and an approximation of the SFDR was determined. In these coding schemes, the accumulator output is truncated to six bits ($S_7, S_6, S_5, S_4, S_3, S_2$). Quarter-wave symmetry is used to drive the fine DAC and half-wave symmetry is used to drive the coarse DAC. Exploiting the symmetry inherent in the sine wave reduces the amount of logic necessary for phase conversion. The four LSBs input to the phase converter ($S_5, S_4, S_3, S_2$) are XORed with $S_6$ to achieve half-wave symmetry for both portions of the DAC. After the conversion logic, the four bits that drive the fine DAC are XORed with the accumulator MSB ($S_7$) to achieve quarter-wave symmetry. The logic that drives the coarse DAC is configured as a thermometer coder, and it represents coarse (or large) steps in the output voltage. The coarse voltage steps are symmetric about the zero-crossing of the sine wave, hence the half-wave symmetry. While the DAC is capable 64 discrete outputs, the phase converter only drives 32 distinct outputs. This is necessary because the DAC steps are all equal size, but the sine step size varies with the sine phase.
Three different coding schemes are investigated. In the first scheme, an ideal DAC that allows for fractions of a unit current is used. This is unrealistic, since it would require a much more complex DAC with more input bits. However, it does provide an upper bound for the maximum SFDR of the DDS, taking the accumulator truncation into account. In the second scheme, the outputs of the DAC are rounded to whole units of current, so that the output levels possible with the DAC that has 3 coarse bits and 4 fine bits are represented. This coding scheme would be used if a ROM were available. With a combinational logic LUT, the logic necessary for implementation is complex. In the third scheme, the LUT from the second scheme is modified to simplify the complicated logic. This reduces propagation delays that could reduce the maximum speed, and it decreases power consumption.

4.5.1.1 Coding with Ideal DAC

The first coding scheme results in the SFDR versus FCW plot shown in Figure 4.20. This scheme has a worst case SFDR of 27.57 dBc at an FCW of 122. Over the whole range of FCWs, the average SFDR is 36.06 dBc. Every fourth FCW, the output SFDR reaches a local maxima, because there is effectively no truncation at these FCWs. This eliminates truncation spurs every fourth FCW. This coding scheme uses an ideal, but unrealistic DAC with a very fine output resolution that is implemented only in MATLAB. The ideal DAC has outputs that have fractions of the unit current weighting. It is useful for providing an estimate of the bounds of the SFDR when using an 8-bit accumulator truncated to 6 bits.

4.5.1.2 Coding with Realistic DAC

A more realistic coding scheme uses DAC output levels that are rounded off to whole values. This scheme provides a simulation of the DAC that is actually implemented in circuitry. This coding scheme would be used if a ROM were available. Since
there is no ROM in this design, combinational logic is used. The combinational logic needed to implement this phase converter is shown in Table 4.2. The SFDR versus FCW plot for this scheme is shown in Figure 4.21. This has a worst case SFDR of 27.57 dBc at an FCW of 122. Over the whole range of FCWs, the average SFDR is 35.69 dBc. As in the first scheme, there are local SFDR maxima every fourth FCW due to the absence of truncation spurs. Since the DAC outputs are rounded off from the ideal case, there is some reduction in the average SFDR. However, the worst case SFDR has the same magnitude and occurs at the same FCW as in the first coding scheme. Although the DAC in this scheme is realistic, the logic needed to implement the LUT shown in Table 4.2 is six levels deep and it requires 38 logic gates.
### Inverted Inputs

\[
A = S_5 \oplus S_6
\]
\[
B = S_4 \oplus S_6
\]
\[
C = S_3 \oplus S_6
\]
\[
D = S_2 \oplus S_6
\]

### Coarse DAC

\[
DAC_6 = S_7
\]
\[
DAC_5 = (A + B \cdot C) \cdot S_7 + \overline{S_7}
\]
\[
DAC_4 = (A + B \cdot C) \cdot \overline{S_7}
\]

### Binary Fine DAC

\[
DAC_3 = ((A \cdot C + A \cdot D) + (B \cdot \overline{C} + (\overline{C} \cdot C) \cdot D) \oplus S_7
\]
\[
DAC_2 = (((A \oplus B) \oplus (C \oplus D)) + (A \cdot B + B \cdot D)) \oplus S_7
\]
\[
DAC_1 = (((A \oplus B) \cdot C + (C \oplus D) \cdot B) + (A \cdot B) \cdot (C \oplus D)) \oplus S_7
\]
\[
DAC_0 = (((A \oplus B) \cdot \overline{C} + (A \oplus B) \cdot \overline{D}) + (A \cdot C) \cdot (B \cdot D))S_7
\]

Table 4.2: Logic necessary to implement the DDSV2 phase converter with a realistic DAC.

![MATLAB simulation of SFDR versus FCW for DDSV2 with coding scheme with the DAC outputs rounded to discrete output levels.](image)

Figure 4.21: MATLAB simulation of SFDR versus FCW for DDSV2 with coding scheme with the DAC outputs rounded to discrete output levels.
4.5.1.3 Simplified Coding with Realistic DAC

If some simplifications are made to the LUT from the second scheme, the complexity and power consumption can be reduced. Using the coding scheme in Table 4.3, the complexity is reduced by one level of depth to five levels of depth, which reduces the LUT propagation delay by one gate delay. Power consumption is also reduced by eliminating 10 logic gates, leaving a total of 28 logic gates in the LUT. Assuming that all logic gates in the LUT use the same amount of power, which is a reasonable assumption for the logic gates in this design, this represents a 26.3% reduction in power by simplifying the coding scheme. The SFDR versus FCW plot for this scheme is shown in Figure 4.22. This has a worst case SFDR of 27.57 dBc at an FCW of 122. Over the whole range of FCWs, the average SFDR is 33.33 dBc. Compared to the other schemes, the SFDR is reduced at many of the FCWs, leading to the reduction in average SFDR over the FCW range. While this scheme has the same worst case SFDR magnitude at the same FCW, some of the FCWs that resulted in local maxima in the previous schemes become local minima in this scheme. In these cases, the simplification added spurs that are not present in the other schemes. This is particularly noticeable in the FCWs that do not have truncation spurs in the previous schemes. The difference is particularly large at FCWs 16, 80, and 96, where the SFDR is 31.97 dBc, 29.53 dBc, and 27.77 dBc respectively. Although there is some loss in performance compared to the second scheme, the savings in complexity and power make the tradeoff worthwhile.

4.5.2 Simulation Results

DDSV2 is constructed by combining the 8-bit accumulator from Section 3.3.2, the phase converter described in Table 4.3, and the DAC from Figure 4.19. To illustrate the maximum clock frequency of the DDS, a simulation output of the DDS is with an FCW of 1 and a 13 GHz clock frequency is shown in Figure 4.23. This simulation shows
Inverted Inputs

\[ A = S_5 \oplus S_6 \]
\[ B = S_4 \oplus S_6 \]
\[ C = S_3 \oplus S_6 \]
\[ D = S_2 \oplus S_6 \]

Coarse DAC

\[ DAC_6 = S_7 \]
\[ DAC_5 = (A + B \cdot C) \cdot S_7 + \overline{S_7} \]
\[ DAC_4 = (A + B \cdot C) \cdot \overline{S_7} \]

Binary Fine DAC

\[ DAC_3 = ((B \oplus C) + A) \oplus S_7 \]
\[ DAC_2 = (\overline{A} \cdot C + B) \oplus S_7 \]
\[ DAC_1 = ((\overline{A} \cdot \overline{B} + \overline{C}) + (\overline{A} \cdot D + B \cdot D)) \oplus S_7 \]
\[ DAC_0 = ((\overline{A} \cdot \overline{B} + D) + (\overline{A} \cdot \overline{C} + A \cdot B)) \oplus S_7 \]

Table 4.3: Simplified logic used in the DDSV2 phase converter to drive the DAC.

![MATLAB simulation of SFDR versus FCW for DDSV2 with simplified coding scheme.](image_url)
Figure 4.23: Simulated time-domain output of DDSV2 with a 13 GHz clock and FCW=1. The plotted output is the difference of the differential signals.

The difference of the differential outputs, so it has a 0 V DC component. At 14 GHz, the DDS has timing glitches and does not operate properly.

The DDS is also simulated over all FCWs to compare the MATLAB simulations of the SFDR to the simulated results. These simulations included extracted parasitic capacitors for the full DDS chip. All 128 FCWs were simulated, taking approximately 64 hours of computer simulation time. The SFDR versus FCW simulation output is shown in Figure 4.24. Compared to the MATLAB simulation of this coding scheme in Figure 4.22, the Cadence Spectre simulation provides a reasonable match. In this simulation, the worst case SFDR is 26.25 dBC, located at an FCW of 96. While this FCW isn’t the minimum in the MATLAB simulation, it is the second to worst FCW. The average SFDR for the simulation of DDSV2 with extracted parasitic capacitors is 33.38 dBC, while it is 33.33 dBC in the MATLAB simulation of the coding scheme.
4.5.3 Measurement Results

The DDSV2 test chip was fabricated as part of the TC9 fabrication run, and it is shown in Figure 4.25. The chip is 2700 µm by 1450 µm and contains 1646 transistors. The circuit was tested on-wafer using similar test setup to the one used for DDSV1, shown in Figure 4.13. Accordingly, time-domain outputs were captured with a sampling oscilloscope and frequency-domain outputs were captured with a spectrum analyzer. The DDS operated up to a maximum clock frequency of 13 GHz, matching the simulation results. The DDS was able to synthesize outputs up to 6.5 GHz in steps of 50.78125 MHz.

With a 13 GHz clock rate and an FCW of 1, the DDS produces the fundamental output frequency of 50.78125 MHz, as shown in the time domain output in Figure 4.26. At this output frequency, the SFDR was measured to be 34 dBc, as shown in the spectrum analyzer output in Figure 4.27. At the maximum FCW of 128 and a 13 GHz clock
Figure 4.25: Microphotograph of the DDSV2 chip. The chip is 2700 \( \mu \text{m} \) by 1450 \( \mu \text{m} \) and contains 1646 transistors.

At this output frequency, the DDS output measured to be 6.5 GHz, as shown in Figure 4.28. At this output frequency, the SFDR is measured to be 50 dBc. The spectrum analyzer output for the 6.5 GHz output is shown in Figure 4.29.

Using automated test software, the SFDR of the DDS was measured over the whole range of FCWs with a 13 GHz clock input. The sweep of SFDR vs. FCW is shown in Figure 4.30. The measured SFDR is better than 30 dBc over most of the FCWs. The worst case SFDR was 26.67 dBc at an output frequency of 6.389775 GHz, which is an FCW of 126. Over the whole range of FCWs, the average SFDR is 33.08 dBc.

Compared to the simulation results, the measured results are a reasonable match. The average SFDR figures are within 0.30 dBc, and the worst case SFDR magnitude is 0.90 dBc lower compared to the MATLAB simulation of the coding scheme, and 0.42 dBc higher compared to the simulation with extracted parasitic capacitors included. While the FCWs with the worst case SFDR are not the same, there is correlation in the shape of the SFDR versus FCW curves, so that many of the local minima and maxima of
Figure 4.26: Measured time-domain output of DDSV2 with a 13 GHz clock and FCW=1. The output frequency is 50.78125 MHz.

Figure 4.27: Measured frequency-domain output of DDSV2 with a 13 GHz clock and FCW=1. The SFDR is 34 dBC.
Figure 4.28: Measured time-domain output of DDSV2 with a 13 GHz clock and FCW=128. The output frequency is 6.5 GHz.

Figure 4.29: Measured frequency-domain output of DDSV2 with a 13 GHz clock and FCW=128. The SFDR is 50 dBc.
SFDR are at the same FCWs. The DDSV2 design consumes 5.42 W of power. DDSV2 is reported in [62].

### 4.6 Summary of Direct Digital Synthesizers

Two DDS designs were presented in this chapter. The DDSV1 design described in Section 4.4 is intended as a high-speed DDS. It uses the ACCV5 8-bit accumulator described in Section 3.3.1, and it also uses a sine-weighted DAC approach. The sine-weighted DAC allows for reduced complexity compared to a traditional ROM lookup table. The DDSV1 design consumes 9.45 W of power and operates up to a maximum clock frequency of 32 GHz. Over the whole range of FCWs, it has a worst case SFDR of 21.56 dBc.

The DDSV2 design described in Section 4.5 is intended as a low power DDS. It uses the low power ACCV6 8-bit accumulator described in Section 3.3.2. At the time of design, a ROM was not available, so DDSV2 instead uses a logic lookup table. Some
simplifications were made to the logic LUT to improve propagation delay and reduce power. These modifications had a small negative impact on the output signal purity, and it is expected that using a ROM would have improved the SFDR. The DDSV2 design consumes 5.42 W of power and operates up to a maximum clock frequency of 13 GHz. Over the whole range of FCWs, it has a worst case SFDR of 26.67 dBc.

A comparison of DDSV1 and DDSV2 to recently reported DDS circuits is shown in Table 4.4. The DDS circuits from this thesis perform better than the recently reported DDS circuits in all areas except SFDR. The speed/power ratio is used as a DDS figure of merit by K. Elliott [42]. DDSV1 has a speed/power ratio of 3.39, which is 5.56 and 2.26 times better than the previously reported designs and 1.41 times better than DDSV2. The DDSV2 speed/power ratio is 3.93 and 1.60 times better than the previously reported results.

While using the speed/power ratio as a figure of merit for DDS designs provides some insight into the relative value of different designs, it has some major flaws. Notably, it does not factor in frequency resolution or SFDR. For example, if a two similar DDS designs are compared, but one has twice the phase resolution, it would need an accumulator twice as large. This would drive up the power consumption, so that even if the two designs had the same maximum clock frequency and SFDR, the design with more frequency resolution would have a worse speed/power ratio. Similarly, if two designs had the same power consumption and the same maximum speed, they would have identical speed/power ratios, even if they had vastly different SFDRs. If all other factors

<table>
<thead>
<tr>
<th>Design</th>
<th>Power (W)</th>
<th>Speed (GHz)</th>
<th>SFDR (dBc)</th>
<th>Speed/Power Ratio (GHz/W)</th>
</tr>
</thead>
<tbody>
<tr>
<td>A. Gutierrez-Aitken et al. [8, 75]</td>
<td>15</td>
<td>9.2</td>
<td>30.00</td>
<td>0.61</td>
</tr>
<tr>
<td>K. Elliott [42]</td>
<td>8</td>
<td>12</td>
<td>30.00</td>
<td>1.50</td>
</tr>
<tr>
<td>DDSV2 [62]</td>
<td>5.42</td>
<td>13</td>
<td>26.67</td>
<td>2.40</td>
</tr>
</tbody>
</table>

Table 4.4: Comparison of recent InP DDS designs.
are equal, a design with a higher SFDR should be consider to be “better” than a design with a lower SFDR, but this is not reflected in the speed/power ratio. Likewise, the extra power consumption necessary to increase the frequency resolution of a DDS should be taken into consideration when comparing designs.

To deal with the shortcomings of the speed/power ratio figure of merit for DDS designs, a new figure of merit is developed that “rewards” improved SFDR and does not penalize for increased frequency resolution. A new figure of merit to consider for DDS circuits is

\[ FOM_{DDS} = \frac{\text{speed} \cdot \text{SFDR} \cdot \text{bitwidth}}{\text{power}}. \] (4.8)

For the high-speed DDS circuits presented in this thesis, the units of \( FOM_{DDS} \) are GHz·dBc-bits/W. The \( FOM_{DDS} \) scales the speed/power ratio up linearly by both the SFDR and the accumulator bit-width. Thus, if all other metrics are equal, a design with twice the SFDR will have twice the \( FOM_{DDS} \). A design with better frequency resolution will also gain a boost in \( FOM_{DDS} \) because the increase in power consumption is be offset by the increase in bit-width.

A calculation of the \( FOM_{DDS} \) for recently reported research and commercial DDS circuits is given for comparison in Table 4.5. The only designs with a better \( FOM_{DDS} \) than DDSV1 and DDSV2 are the AD9858 \[41\] part, which does not report a full-Nyquist SFDR, so it is not an accurate comparison of \( FOM_{DDS} \), and the low power DDS designs \[46,47\]. DDSV1 and DDSV2 have a better \( FOM_{DDS} \) than the InP DDS state of the art designs. It should be noted that the DSSV1 design is penalized in \( FOM_{DDS} \) in comparison to the speed/power ratio, since it has a lower SFDR than the other designs. Compared to the InP designs only, the \( FOM_{DDS} \) for the DSSV1 design is 3.97 times that of the A. Gutierrez-Aitken et al. \[8\] design, 1.62 times the K. Elliott \[42\] design and 1.14 times the DSSV2 design. Since the DSSV2 design has a similar SFDR to the previously reported designs, it is not penalized as much, and its
<table>
<thead>
<tr>
<th>Part Name</th>
<th>$f_{clk}$ (GHz)</th>
<th>FCW (bits)</th>
<th>SFDR (dBc)</th>
<th>Power (W)</th>
<th>$FOM_{DDS}$ (GHz·dBc·bits/W)</th>
<th>Ref.</th>
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<td>107</td>
<td>[43]</td>
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<td>None$^d$</td>
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<td>1.6</td>
<td>30</td>
<td>40$^b$</td>
<td>6</td>
<td>320</td>
<td>[45]</td>
</tr>
<tr>
<td>None$^d$</td>
<td>12</td>
<td>10</td>
<td>30</td>
<td>8</td>
<td>450</td>
<td>[42]</td>
</tr>
<tr>
<td>DDSV2$^d$</td>
<td>13</td>
<td>8</td>
<td>26.67</td>
<td>5.42</td>
<td>512</td>
<td>[62]</td>
</tr>
<tr>
<td>DDSV1$^d$</td>
<td>32</td>
<td>8</td>
<td>21.56</td>
<td>9.45</td>
<td>584</td>
<td>[51]</td>
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<tr>
<td>None$^d$</td>
<td>2.0</td>
<td>8</td>
<td>35</td>
<td>0.82</td>
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<td>[47]</td>
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<tr>
<td>AD9858</td>
<td>1.0</td>
<td>32</td>
<td>50$^c$</td>
<td>2</td>
<td>800</td>
<td>[41]</td>
</tr>
<tr>
<td>None$^d$</td>
<td>0.8</td>
<td>32</td>
<td>25</td>
<td>0.174</td>
<td>3678</td>
<td>[46]</td>
</tr>
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</table>

$^a$ Frequency Control Word (FCW) or accumulator resolution.

$^b$ Claims 20 dBc spectral purity for harmonics.

$^c$ For 360 MHz output.

$^d$ Uses InP DHBT Technology.

Table 4.5: Recent commercial and reported direct digital synthesizers compared using $FOM_{DDS}$.

$FOM_{DDS}$ is 3.48 times the A. Gutierrez-Aitken et al. [8] design and 1.42 times the K. Elliott [42] design.

The new $FOM_{DDS}$ figure of merit from Equation 4.8 should prove useful for comparing DDS designs. It includes more information on the metrics that are important for DDS performance, namely SFDR, frequency resolution, clock frequency, and power consumption, and improves upon previous DDS figures of merit.
CHAPTER 5

Conclusion

In this thesis, high-speed digital and mixed-signal circuits, with clock frequencies in the range of 12 GHz to 41 GHz, have been implemented in the Vitesse VIP-2 InP DHBT technology [1]. The main focus of the work was on maximizing the performance of accumulator and DDS circuits that are suitable for radars and communications systems in the X-band and Ku-band range of 8 GHz to 16 GHz, while minimizing the power consumption. The two conflicting tradeoffs of high-speed and low power lead to several accumulator and DDS designs over a range of speed/power design points.

This thesis also focused on issues with the simulation of transmission lines for on-chip clock signals, that become an issue for clock frequencies above 30 GHz routed on interconnects with lengths above 300 \( \mu \)m. At such high frequencies and long interconnect lengths, the clock distribution lines can no longer be modelled as lumped element lines, and instead must be dealt with as transmission lines. These transmission line models are not a direct part of the design kit, and must be handled in a special way in order to properly capture their behavior.

Finally, a new and improved figure of merit for comparing DDS designs was developed and used to compare DDS designs. This figure of merit incorporates more of the important metrics of DDS performance, such as SFDR and frequency resolution, than previous figures of merit, so it provides a more complete method of design comparison.

The major accomplishments of this thesis and recommendations for future work are summarized below.

5.1 Summary of Accomplishments

The first major accomplishment of this thesis was the development of the ACCV2 accumulator described in Section 3.2.4. This accumulator was fabricated as part of the
<table>
<thead>
<tr>
<th>Fabrication Run</th>
<th>Date</th>
<th>Designs</th>
</tr>
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<tr>
<td>TC6</td>
<td>January 2004</td>
<td>ACCV1, ACCV2</td>
</tr>
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<td>TC7</td>
<td>June 2004</td>
<td>ACCV1, ACCV3, ACCV4</td>
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<td>TC8</td>
<td>November 2004</td>
<td>ACCV5, DDSV1</td>
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<td>TC9</td>
<td>March 2005</td>
<td>ACCV6, DDSV2</td>
</tr>
</tbody>
</table>

Table 5.1: Summary of fabrication runs, dates, and designs in each fabrication run.

TC6 fabrication run. Results of its performance were published in 2005 [50]. At the time of its publication, it was the fastest reported 4-bit accumulator, with a maximum clock frequency of 41 GHz and a power consumption of 3.04 W. The previous state of the art accumulator circuit reported by T. Mathew et al. [27, 59] in 2001 only consisted of accumulator components, and only operated up to a maximum clock frequency of 19 GHz. Thus, not only was ACCV2 more than twice as fast as the previous state of the art, but it was also integrated as a complete accumulator, instead of just accumulator components. The ACCV2 accumulator included a novel single-level parallel-gated carry circuit design. This was the first time that this carry circuit had been demonstrated. This carry circuit enabled reduced power consumption in later accumulator versions. In the next generation fabrication run TC7, ACCV3, which is described in Section 3.2.5, used the single-level parallel-gated carry circuit along with a three-level series-gated sum circuit to improve the speed/power ratio compared to ACCV2 from 13.49 GHz/W to 17.26 GHz/W. Table 5.1 summarizes the fabrication runs, dates, and the designs in each fabrication run.

Further iterations of accumulator designs extended the bit-width to 8-bits for use in DDS circuits. Two 8-bit accumulator circuits described in Section 3.3 show further improvement in the area of speed/power ratio. For comparison, if the ACCV3 4-bit accumulator is extended to 8-bits, it would have a speed/power ratio of 5.84 GHz/W. The ACCV5 that was part of fabrication run TC8 operated up to a 32 GHz clock frequency and has a speed/power ratio of 6.54 GHz/W. In fabrication run TC9, the low power ACCV6 operated up to 13 GHz and had a speed/power ratio of 6.10 GHz/W. Both
designs showed improvement over the ACCV3 design. It should also be noted that although the ACCV6 design had a lower speed/power ratio than ACCV5, it would have a better speed/power ratio if the designs were extended to larger bit-widths, because it uses fewer registers, and the registers will become the dominant portion of power consumption when the bit-width is extended.

As part of the implementation of high-speed digital and mixed-signal circuits, it became necessary to develop a method for dealing with transmission line effects for interconnects. This was of particular importance for the clock signal interconnects, since these interconnects operate at the highest frequency on the chip and tend to have the longest lengths, since they must route clock signals to many locations on the chip. Since the design kit did not automatically include the transmission line models, it was necessary to create separate models of the clock tree with the line parameters and loads used in layout, as described in Section 1.1.2. These models were then simulated with AC simulation in the Cadence Spectre circuit simulator to determine the frequencies over the area of interest that had excessive gain or attenuation. The gain or attenuation could then be pinpointed and dealt with by adding series resistance or clock buffers. Once the AC simulations showed reasonable results, there was sufficient confidence in the clock tree for use in transient simulations, where the transmission line effects are not modelled. Since the high-speed clock rates and interconnection lengths are approaching transmission lines, and the simulation environment does not directly handle these effects, this method of handling the transmission lines was necessary. This approach had not been previously discussed in literature. In some design kits for other processes, microwave models are included for the transmission lines, however, that is not the case in the Vitesse VIP-2 process. It is expected that the transmission line effects will only become more of a problem and begin to affect signal interconnects as well as clock interconnects as clock frequencies increase in the near future and the interconnects become
electrically long as defined by Equation [1.1]. To deal with this problem, either design kits will need to address the issue, or the method outlined above can be utilized.

The second major accomplishment of this thesis was the design of high-speed DDS circuits that resulted in the fastest reported DDS circuits. Previously, the fastest DDS circuits were reported by A. Gutierrez-Aitken et al. [8] in 2001, operating up to a 9.2 GHz clock frequency, and by K. Elliott [42] in 2005, operating up to a 12 GHz clock frequency. DDSV1, the first generation DDS, which was described in Section 4.4, operated up to a maximum clock frequency of 32 GHz and consumed 9.45 W of power. This greatly exceeded both of the previous designs. The second generation DDS named DDSV2 was described in Section 4.5. Even though DDSV2 was designed as a low power DDS, its performance and power consumption exceeded that of the A. Gutierrez-Aitken et al. [8] and K. Elliott [42] designs. It operated up to a maximum clock frequency of 13 GHz with a 5.42 W power consumption.

There was also improvement over the previous state of the art in terms of the speed/power figure of merit presented by K. Elliott [42]. The A. Gutierrez-Aitken et al. [8] DDS had a speed/power ratio of 0.61 GHz/W and the K. Elliott [42] DDS had a speed/power ratio of 1.50 GHz/W. In comparison, DDSV1 has a speed/power ratio of 3.39 GHz/W and DDSV2 has a speed/power ratio of 2.40 GHz/W.

While the speed/power figure of merit is often used and provides some idea of relative accumulator performance, it only factors in the speed and power of a DDS. Thus, it omits important DDS metrics such as SFDR and frequency resolution. To provide a better means of comparing DDS circuits, a new figure of merit for DDS circuits ($FOM_{DDS}$) was devised in this thesis and defined in Equation 4.8. This metric modifies the speed/power figure of merit used by K. Elliott [42] to include SFDR and frequency resolution (through the accumulator bit-width). This new figure of merit more completely captures the important metrics of DDS, so it is better suited for comparing such
circuits, and a comparison to recently reported DDS circuits, in InP and other technologies, was shown in Table 4.5. The only designs with a better $FOM_{DDS}$ than DDSV1 and DDSV2 are the AD9858 [41] by Analog Devices, which does not report a full-Nyquist SFDR, so it is not an accurate comparison of $FOM_{DDS}$, and the low power DDS designs by B.-D. Yang et al. and X. Yu et al. [46, 47]. Compared to the previous state of the art InP designs by A. Gutierrez-Aitken et al. [8] and K. Elliott [42], DDSV1 [51] and DDSV2 [62] have an improved $FOM_{DDS}$.

5.2 Recommendations for Future Work

Based on the work in this thesis, three recommendations for future work are identified: DDS circuit improvements for higher SFDR, implementation of similar circuits in SiGe technology, and the improvement of the interconnect technology along with the development of microwave models for interconnects.

The first recommendation is to push for DDS circuits with improved SFDR. This would be accomplished by improving all of the blocks of the circuit. First, the bit-width of the accumulator could be extended to increase the frequency resolution of the DDS. This would also provide a larger phase bit-width for the phase conversion logic. The accumulators in this thesis have an architecture that facilitates extensions in bit-width, but there would likely be some challenges from the increased clock distribution and power consumption.

The next area of the DDS to improve would be the phase conversion circuitry. In this thesis, DDSV1, which is described in Section 4.4, used a phase conversion that was integrated with the DAC, and DDSV2, which is described in Section 4.5, used a logic LUT. These phase conversion schemes were fairly simple, but they were not very flexible. They also only allowed for a worst case SFDR that was below 30 dBc. Using a ROM would allow for more flexibility, in that the ROM coding can be changed much easier than a logic coding. A 16 by 6 bit ROM with a 36 GHz clock frequency has
been reported S. Manandhar et al. [78] using this InP technology. The 36 GHz clock frequency leaves margin for expanding the size of the ROM and maintaining high-speed performance so that it could be integrated into either of the DDS designs. Using the ACCV5 circuit with a ROM as a phase converter and a different DAC circuit would be likely to lead to a DDS with an improved SFDR, with a worst case SFDR closer to 30 dBC, while maintaining the 32 GHz clock frequency. This would also lead to an improvement in the $FOM_{DDS}$. The ROM would be particularly advantageous in the DDSV2 accumulator. As noted in Section 4.5.1, the logic LUT coding scheme was simplified from what could have been used had a ROM been available. Although the simulation results did not show a difference in the worst case SFDR, the average SFDR was simulated to be 33.33 dBC with the simplified logic LUT and 35.69 dBC with a ROM. A ROM would not have to be used alone, either. A combination of the ROM with some logic could allow for a coding scheme that would provide improvements SFDR beyond 40 dBC.

Finally, more work could be done on DAC design. In order to improve the SFDR, not only would the phase converter need to be improved, but the DAC would need an extended bit-width. There is much work to be done in the area of high-speed DAC design, but it is essential for new high-speed DDS circuits with worst case SFDR above 40 dBC. By further investigating all of the components of the DDS, the DDS operation could be improved and allow for more applications.

The second recommendation for future work is to duplicate and extend this work in SiGe. InP was used in this thesis, because one of the aims of the TFAST project was to push the development of InP technologies. There is interest in developing InP technologies because their superior $BV_{CEO}$ compared to SiGe is expected to result in high performance mixed-signal circuits. Despite the lower $BV_{CEO}$, SiGe technology has some advantages in other areas that could be exploited. First of all, since the SiGe process is built on top of a CMOS process, PMOS and NMOS transistors are available.
An NMOS transistor could be used as a high impedance current source with a lower voltage drop than the current sources in this thesis. The high impedance of the NMOS source would be an improvement over the low impedance resistor-only current sources that had a low voltage drop. This would allow for a reduction in power consumption. Secondly, the SiGe process has copper interconnects. The InP process is currently limited to aluminum interconnects with wide pitches compared to the copper interconnects in SiGe. Smaller interconnects would allow for reduced parasitic capacitances and improve the relative performance. Finally, the SiGe process is more manufacturable than InP, so yields and consistency between fabrication runs would be improved. Since SiGe is more manufacturable than InP, it also costs less. SiGe is rapidly developing due to interest in RF technologies, so at some point in the future, the performance of SiGe devices may exceed the performance of InP devices. Some of the designs in this thesis could be ported to a SiGe technology, to compare the mixed-signal performance of SiGe vs. InP. The designs could be modified for reduced power by using NMOS current sources, and the differences in layout parasitic capacitance could be explored.

The final recommendation for future work is an improvement of the interconnect technology and the development of microwave models for interconnects. As described in Section 2.1, the Vitesse VIP-2 technology has four levels of aluminum interconnects. The minimum size of the aluminum interconnect is 1 µm, so parasitic capacitances quickly add up and hamper high-speed performance. The interconnect is also not planarized, so the extraction of parasitic capacitances is prone to errors. Moving to a planarized copper interconnect technology with finer feature sizes would allow for a reduction in parasitic capacitances that would lead to improved performance. This would also neutralize one of the advantages that SiGe technologies currently have over this InP technology.

With planarized interconnects, microwave models for the interconnects could also be developed. These would be useful for the clock tree or any other signal lines that
are electrically long by Equation 1.1. The microwave models would greatly simplify simulation by eliminating the extra steps that were developed in this thesis and described in Section 2.3. This too would eliminate one of the advantages that some other toolkits have over this InP technology.
REFERENCES


BIOGRAPHY OF THE AUTHOR

Steven Turner was born in Portland, Maine. He received his high school education from Westbrook High School in Westbrook, Maine in 1997.

He entered The University of Maine in 1997 and obtained his Bachelor of Science degree in Electrical Engineering and Computer Engineering in May 2001 and his Master of Science degree in Computer Engineering in May 2003.

Since September 2001, he has served as a Research Assistant at The University of Maine. During the summer months he has served internships at Tundra Semiconductor in South Portland, Maine and BAE Systems in Nashua, New Hampshire. His current research interests include high-speed digital and mixed-signal microelectronics design. He is a member of IEEE, Tau Beta Pi, and Eta Kappa Nu.

He is a candidate for the Doctor of Philosophy degree in Electrical Engineering from The University of Maine in May 2006.